

Very-efficient flipping of many coins

(between two parties)



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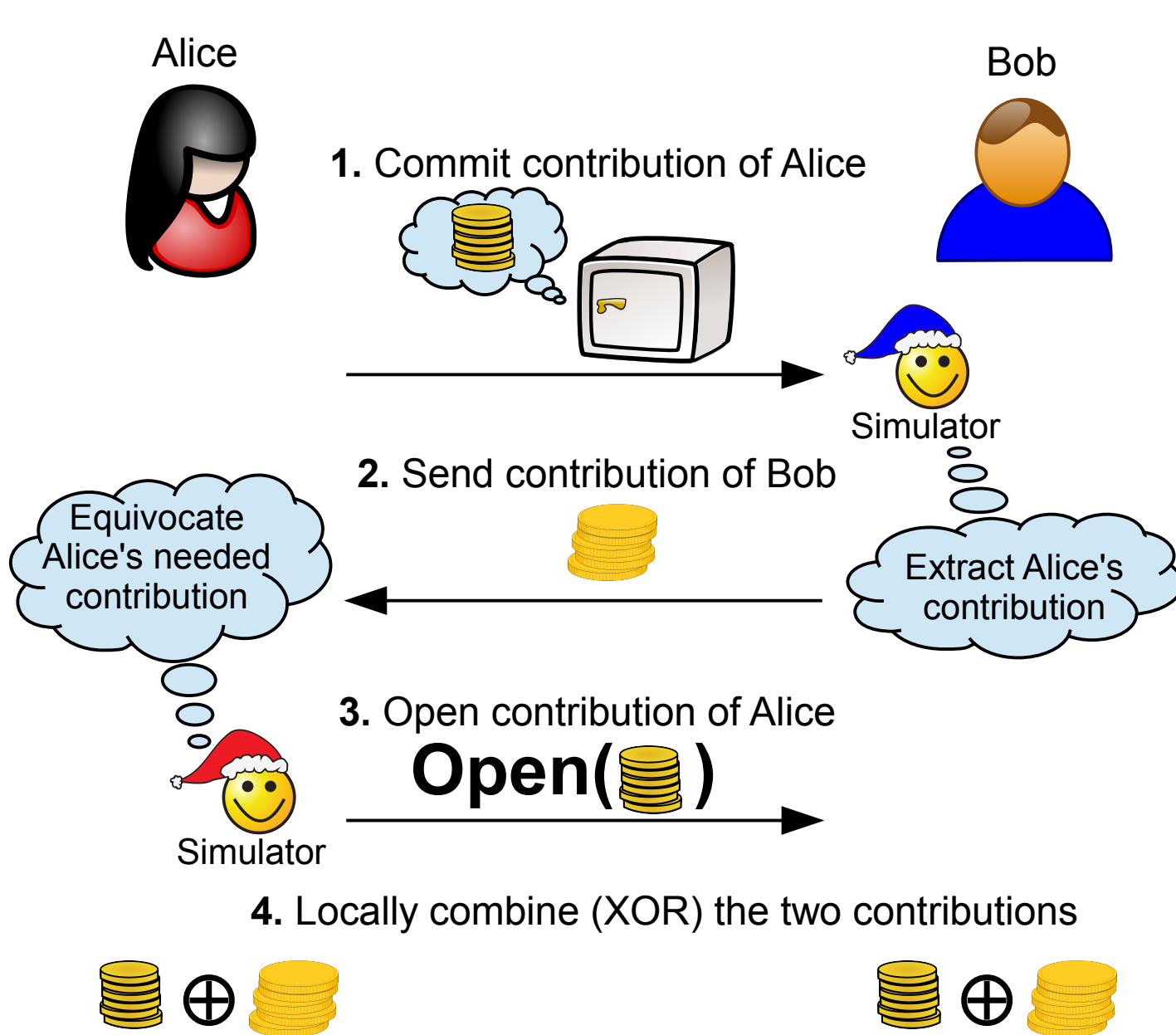
University of Lisbon and Carnegie Mellon University

Early presentation of results at *rump session*
of Theory of Cryptography Conference 2014
February 25, San Diego, USA

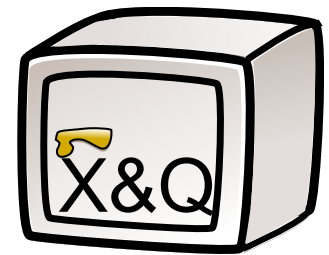
(Minor adjustments on Feb 28, when preparing upload to Internet)

The author is a Ph.D. student at FCUL-DI and CMU-ECE. Support for this research was provided by the Fundação para a Ciência e a Tecnologia (Portuguese Foundation for Science and Technology) through the Carnegie Mellon Portugal Program under Grant SFRH/BD/33770/2009.

The traditional coin-flipping template

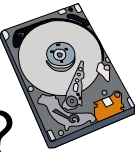


Commitment scheme needs be both **extractable (X)** and **equivocable (Q)**, i.e., be X&Q.



Several constructions exist ... with group-elements or group-operations in number or with size proportional to # coins

(Two recent independent works devise more efficient methods)

What if we want to flip **MANY** coins, e.g., 2TB? 

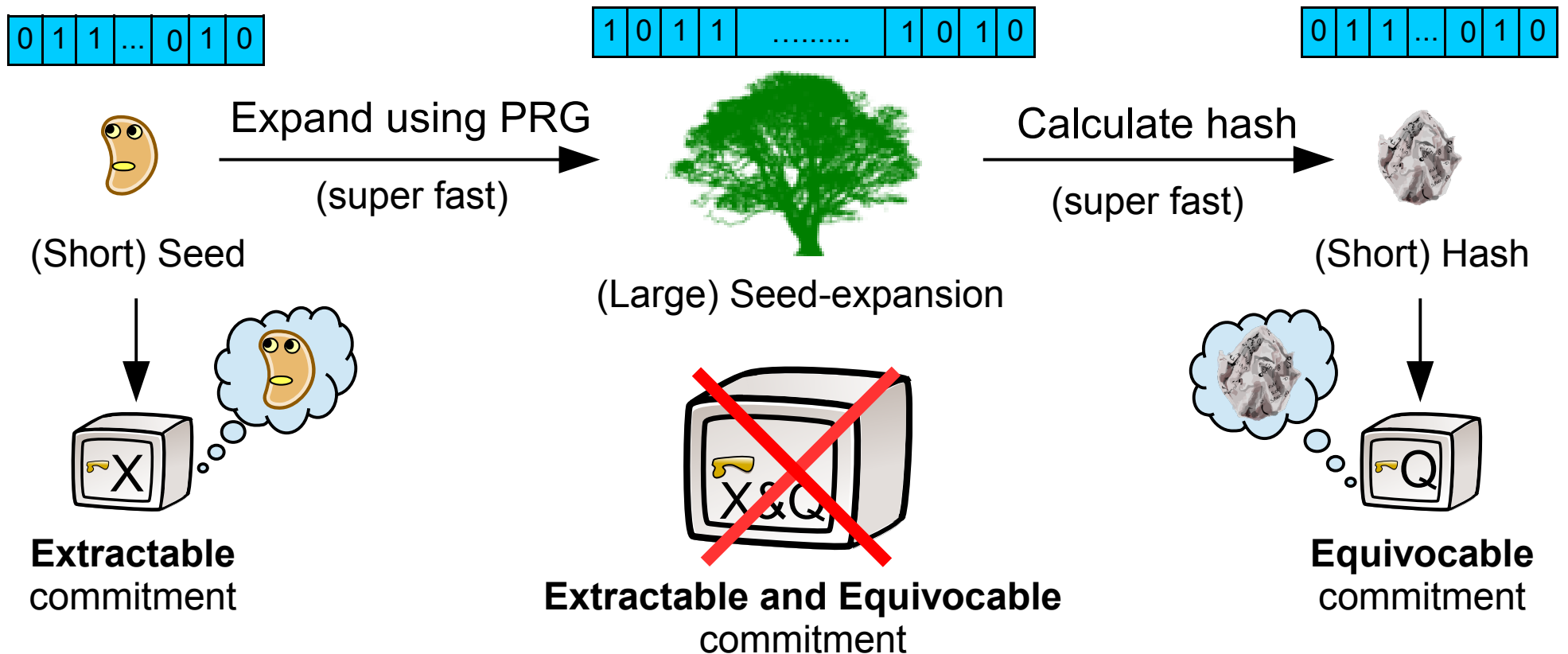
A new approach

- Can we achieve a BIG X&Q commitment using only:
- a FEW SMALL X-commits and a FEW SMALL Q-commits;
 - and symmetric primitives (PRG, hash function, XORs)

?

Yes!

An initial intuition

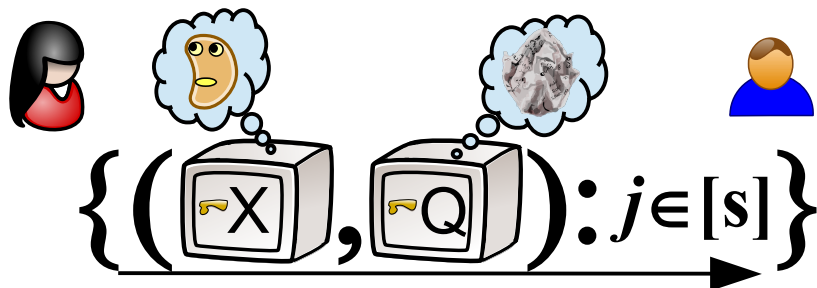


One-pass simulatable coin-flipping

(cut-and-choose based technique)

0. Prepare seeds and hashes

0.1. Alice commits seeds and hashes



0.2. Cut-and-choose: $[s] = J_V + J_E$

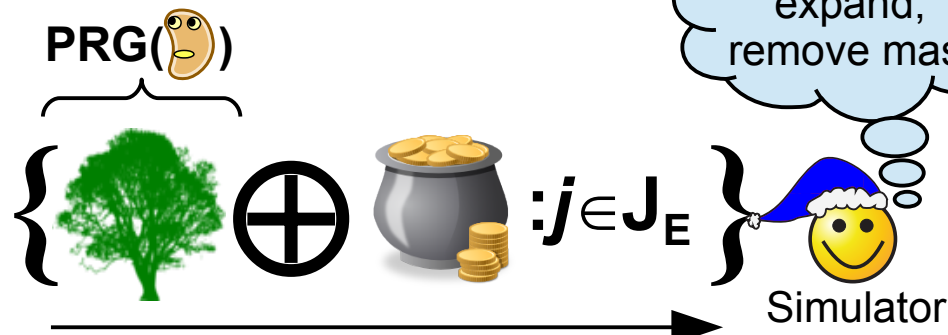
0.3. Alice opens *verification* instances (J_V)



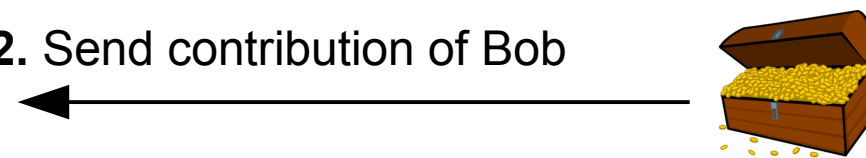
0.4. Bob verifies (J_V): $\text{hash} = \text{Hash}(\text{PRG}(\text{seed}))$

\Rightarrow high probability that a portion of remaining instances (J_E) are consistent

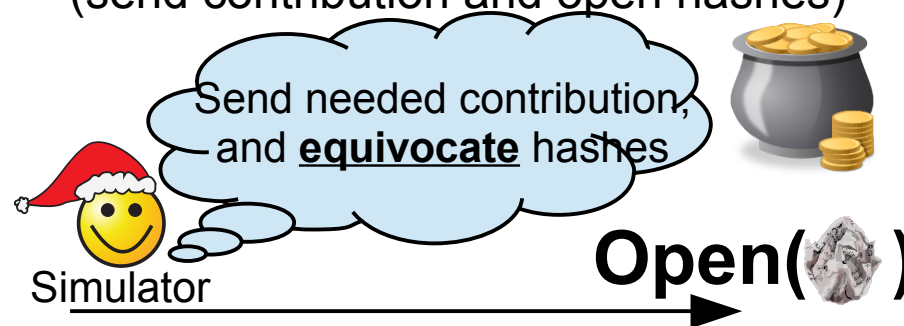
1. X-Commit contribution of Alice



2. Send contribution of Bob



3. Q-open contribution of Alice (send contribution and open hashes)



4. Locally combine contributions:



Summary

- A new approach for flipping many coins
 - Uses few X -commits of seeds and Q -commits of hashes
 - Leverages throughput of PRG and hash function
- Overlooked in this short presentation:
 - Verifiability condition for simulator to check that extracted hash is consistent with masked contribution.
 - How to reduce communication, by fragmenting the contribution using an efficient *Information dispersal algorithm* (and respective reconstruction).
 - Probabilities associated with the cut-and-choose.
 - A much simpler solution exists if rewinding is allowed.

Thank you for your attention!

Very-efficient flipping of many coins

(soon to be on eprint)

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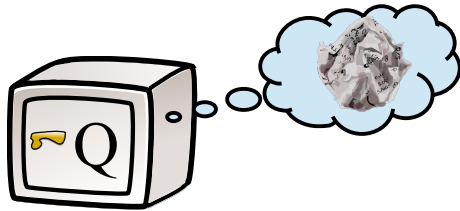
The images used in this presentation were taken from openclipart.org and clker.com, with the expectation of being in the public-domain and available for free usage.

If rewinding is OK, use another template

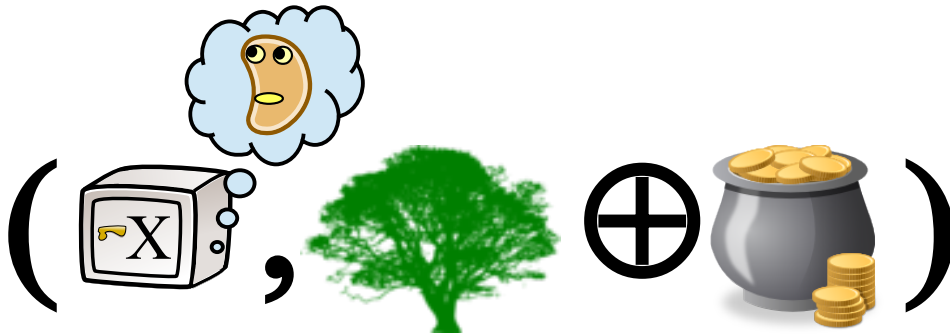
(Slide prepared for rump session but not shown due to time constraint)



1. Bob Q-commits hash of his contribution



2. Alice X-commits seed of a mask, and sends her masked contribution



PRG()



3. Bob sends his contribution and Q-opens its hash

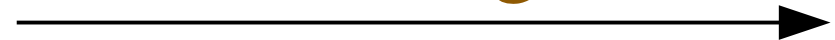


Open()



4. Alice opens the seed of her contribution mask

Open()



5. Locally unmask contribution of Alice and combine contributions

