The background of the slide is a grayscale image of a printed circuit board (PCB) with various components and traces. A horizontal blue line with a small step-down on the left side is positioned below the title.

Breaking Ciphers with COPACOBANA A Cost-Optimized Parallel Code Breaker or How to Break DES for 8,980 €

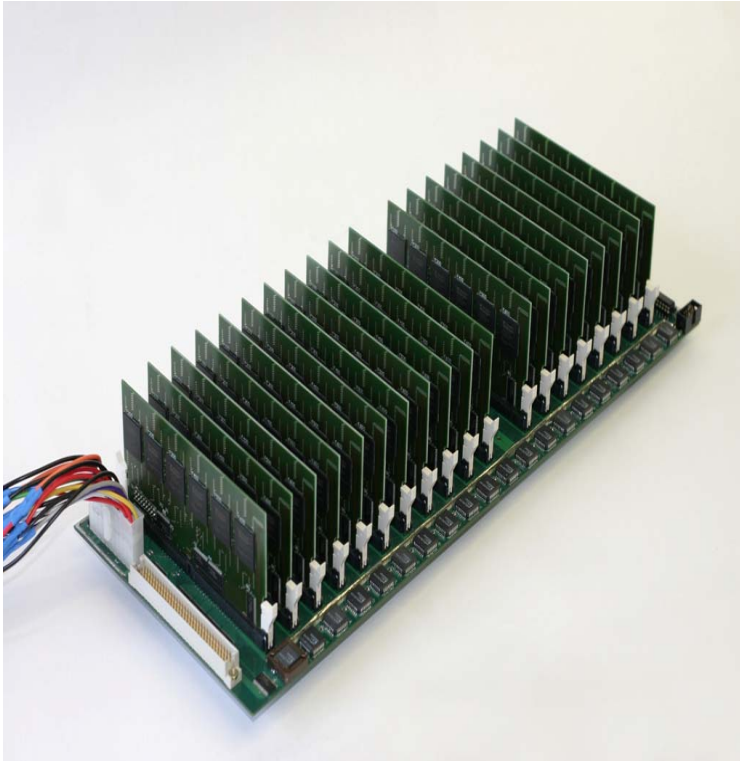
CHES 2006, Yokohama, October 10-13, 2006

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Manfred Schimmler, Christof Paar

Acknowledgements

- Joint project with the University of Kiel (Gerd Pfeiffer, Manfred Schimmler)
- Special thanks to François-Xavier Standaert and Jean-Jacques Quisquater (Université Catholique de Louvain) for the core of the DES architecture

What's in a name?



Copac**o**bana



- **Security vs. Cost**
- COPACOBANA Design
- Application 1: Brute Force Attack on DES
- Application 2: ECC Attack
- Conclusion and Outlook

When is a Cipher Secure?

Symmetric ciphers

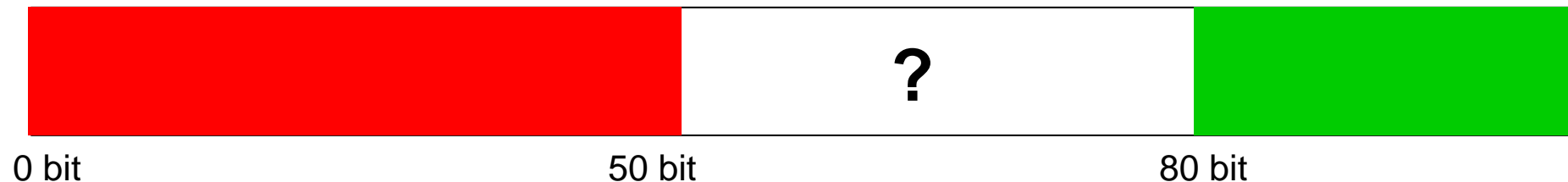
- (hopefully) only brute-force attack possible
- „secure“ key lengths: 112...256 bit (attack compl. $2^{112} \dots 2^{256}$)
- but in practice wide variety of keys: AES, DES, RC4, A5, MD5, SHA-1, ... (attack compl. $2^{56} \dots 2^{256}$)

Asymmetric ciphers (RSA, ECC, DL)

- algorithmic attacks (e.g., factorization) dictate larger keys
- key lengths in practice:
 - RSA, DL: 1024 ... 4096 bit
 - ECC: 160 ... 256 bit
- attack complexities: 2^{80} (?) ... 2^{128}

Security and Computation

- Traditional: security of ciphers = **complexity** of attacks
- However: What really matters are the **costs** of an attack
- State-of-the-art
 - $< 2^{50}$ steps can be done with PC networks (more or less conveniently)
 - $> 2^{80}$ steps are very hard with today's technology (probably also for intelligence agencies)



Major question: Cost of attack for ciphers with 50...80 bit security
(RSA1024, ECC160, SHA-1, DES, A5, ...)

Introduction: Massive Computing

Supercomputers (Cray, SG, ...)

- General (= complex & expensive) parallel computing architectures
- fast I/O, large memory, easy to program
- ▶ **poor cost-performance ratio for (most) cryptanalysis**



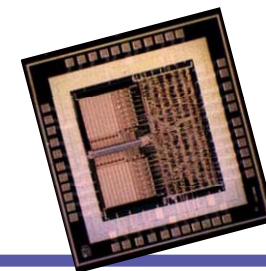
Distributed computing (conventional PCs)

- Dedicated clients in clusters, or
- Using PC's idle time: E.g., SETI@home (BOINC framework)
- ▶ **Problem of motivation, confidentiality issues**



Special-purpose hardware

- ASIC - Application Specific Integrated Circuits (high NRE)
- FPGA - Field Programmable Gate Arrays (low NRE)
- ▶ **best cost-performance ratio**



Introduction: Advantage of Hardware

Cost-performance ratio of DES¹⁾: PC vs. FPGA

- DES encryptions / decryptions per second



Pentium4@3GHz: $\approx 2 \times 10^6$
price per device (retail): € 80



Xilinx XC3S1000@100MHz $\approx 400 \times 10^6$
price per device (retail): € 40

► Cost-performance ratio differs by 2-3 orders of magnitude!

1) Based on actual optimized implementations

COPACOBANA: Design Principles

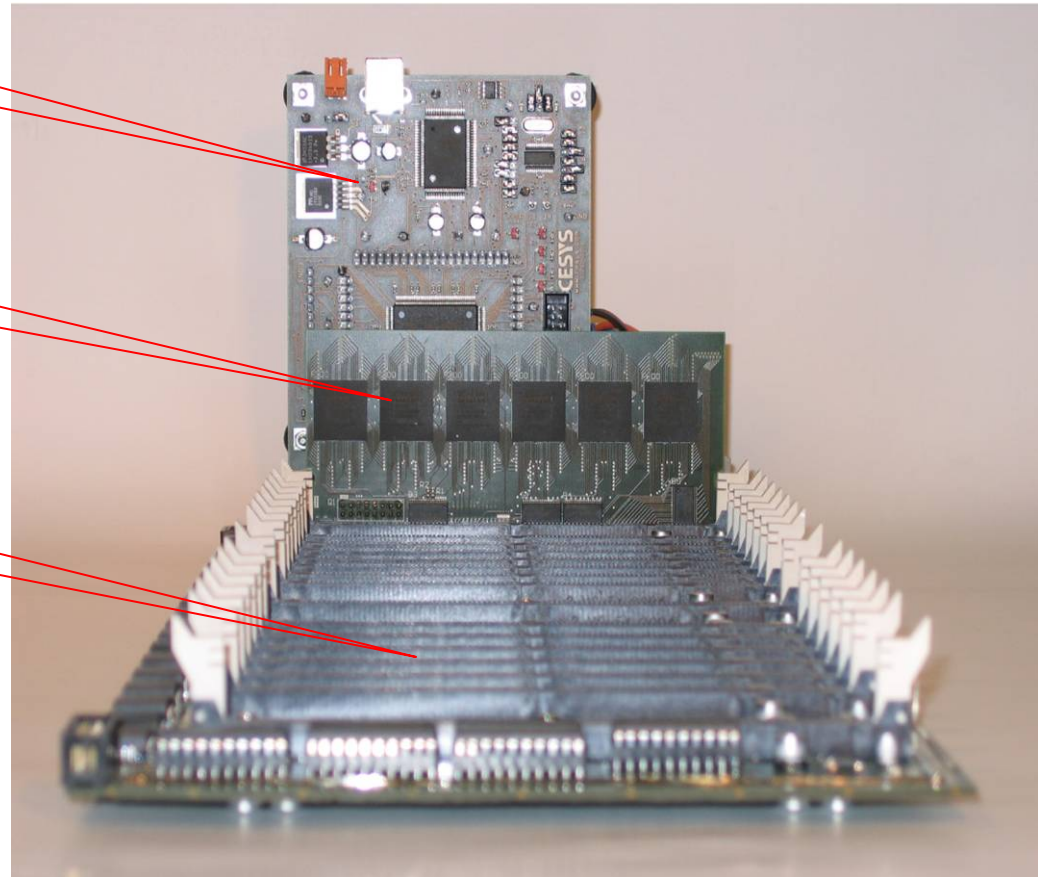
- Ability to perform $\geq 2^{56}$ crypto operations
- **Re-programmable**: Applicable to many ciphers
- Strictly optimized **cost-performance ratio**:
 - off-the-shelf hardware (low-cost)
 - many logic resources (performance)
- **< 9,000 €** (including fabrication and material cost)
- **Parallel** architecture, based on 120 low-cost FPGAs
- Sacrifices
 - no global memory
 - no high-speed communication („only“ Mbit/s)

COPACOBANA: Realization

Controller
board

FPGA module

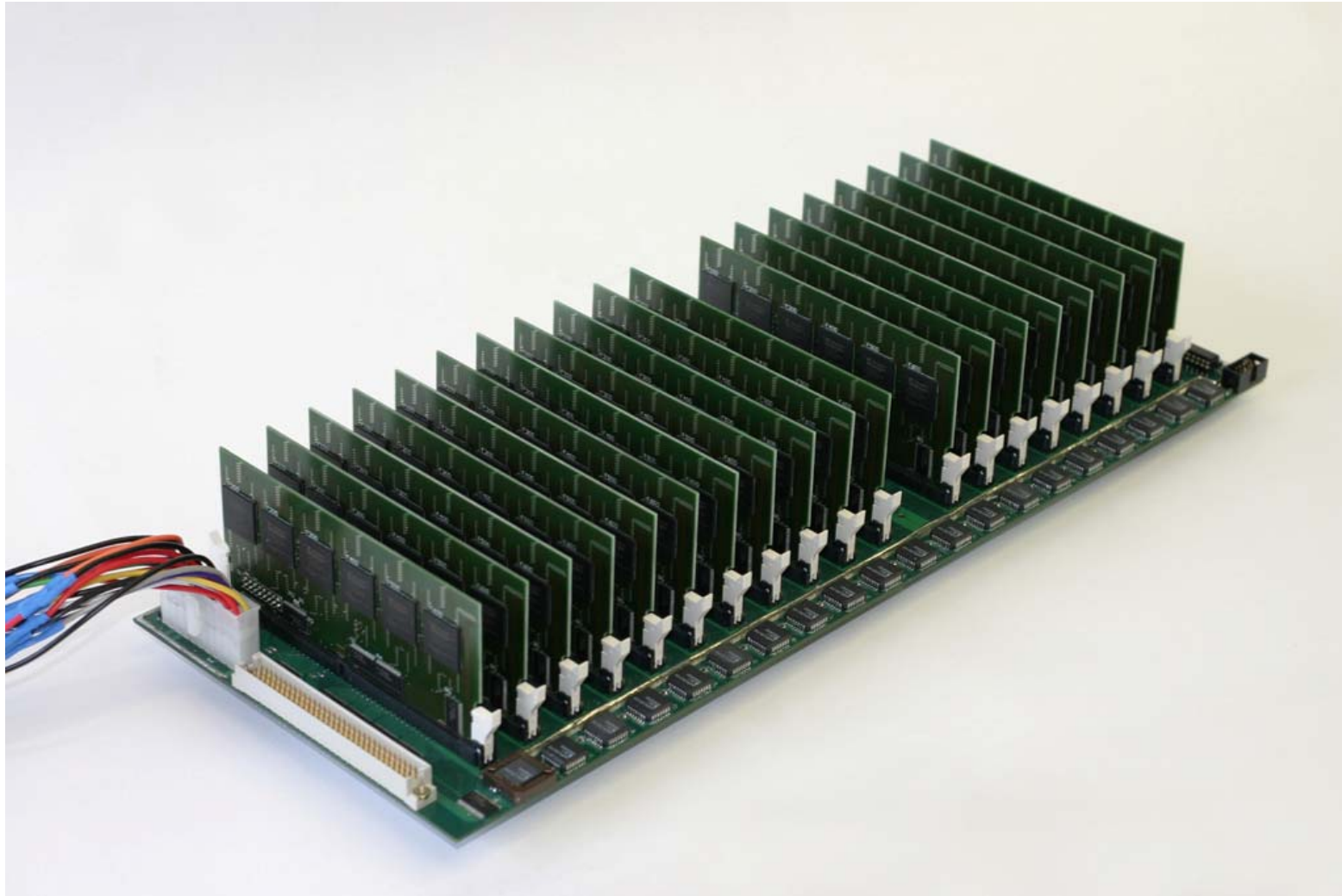
Backplane



Scales easily:

- 20 FPGA modules/machine
(120 FPGAs/machine)
- multiple machines via USB

COPACOBANA: Alpha Prototype



COPACOBANA: Applications

First flexible cryptanalytical machine outside government agencies

1. Exhaustive key search of DES

- ciphers with $2^{56} \dots 2^{64}$ attack steps possible

Attacks
feasible

2. Real-world systems such as ePass, Norton Diskreet, ...

Robust
security
estimations

3. Elliptic Curve Discrete Logarithm Problem (ECDLP)

- Parallelized Pollard's Rho

Improves other
attacks

4. Factorization

- Parallelized Elliptic Curve Method (ECM) as subroutine for GNFS (see GMU's talk later)

Outline

- Security vs. Cost
- COPACOBANA Design
- **Application 1: Brute Force Attack on DES**
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Cryptanalytical Applications: Attacks on DES

Data Encryption Standard (DES):

- Block cipher with 56-bit key
- Expired standard, but still used (legacy products, ePass, Norton Diskreet, ...)

Exhaustive key search (conventional technology):

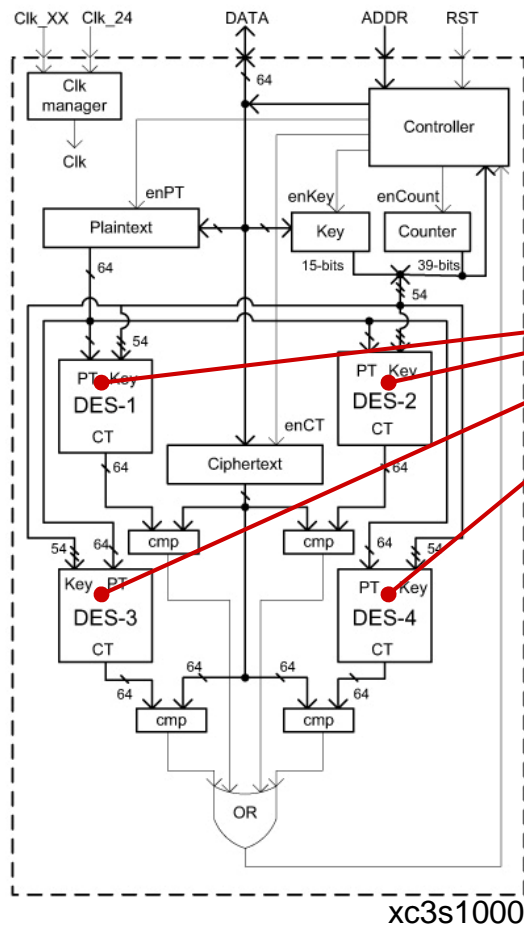
- Check 2^{55} keys on average
- PC (e.g., Pentium4@3GHz) \approx 2 mio. keys/sec
- Average key search with one PC $\approx 2^{34}$ sec = 545 years!



► Can do much better with special-purpose hardware!

Attacks on DES

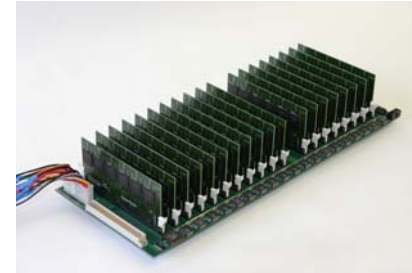
FPGA-based attacks on the Data Encryption Standard (DES):



- Exhaustive key search (FPGA based):
 - 4 completely pipelined DES engines per FPGA (courtesy of the crypto group of UCL)
 - one key per clock cycle per DES engine
 - One FPGA @ 100MHz: 400 mio. keys/ sec

Attacks on DES

- COPACOBANA: average key search of **8.7 days** @ 100 MHz
- Somewhat higher clock rates possible
- FPGA vs. PC (average key search in 8.7 days)
 - 22,865 Pentium 4 (€ 3.6 million incl. overhead)
 - or
 - COPACOBANA (total cost € 9000 incl. overhead)
- Alpha version of COPACABANA runs stable
- Life attack at <http://www.copacobana.org/live>



A Historical Perspective: The Power of Moore's Law

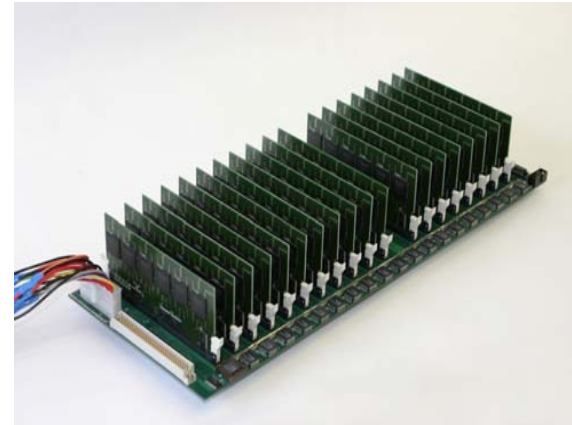
DeepCrack, 1998

\$250,000



COPACOBANA, 2006

\$10,000



Moore's Law: 50% cost reduction / 1.5 years

2006-1998 = 8 years $\approx 5 \times 1.5$ years

Prediction: $\$250,000 / 2^5 \approx \$8,000$ (close to actual \$10,000)

Outline

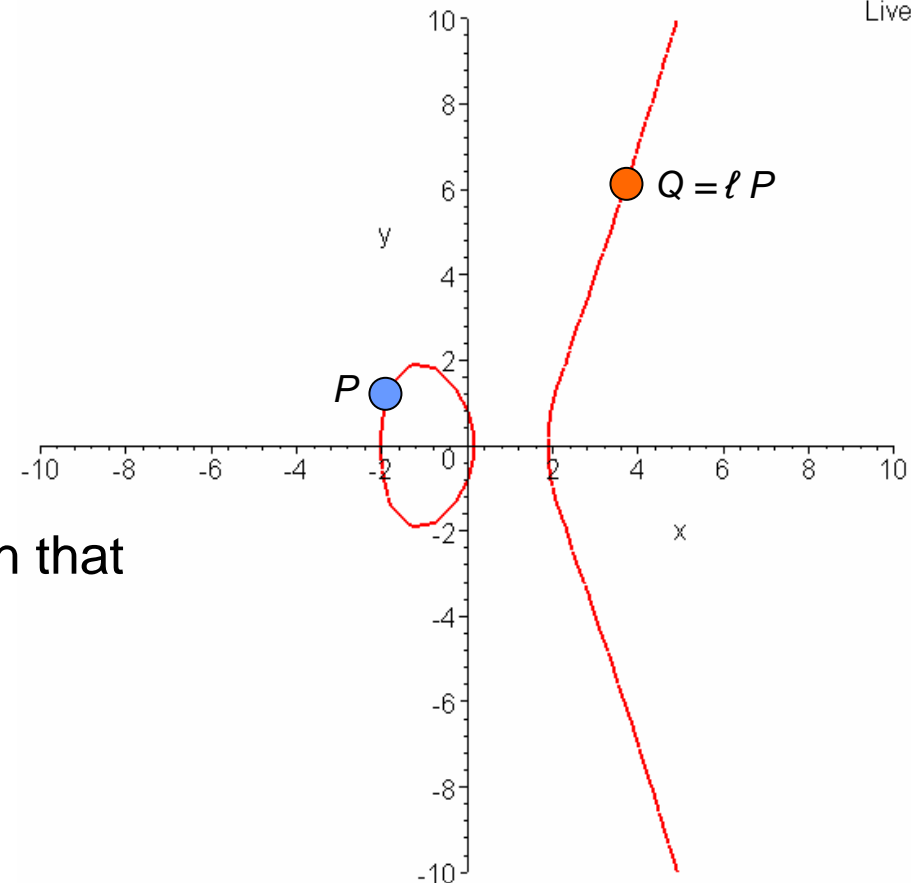
- Security vs. Cost
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ECDL Problem

Live

- Many real-world applications rely on hardness of ECDLP
 - ECDSA,
 - ECDH,
 - ...
- Let P be a generator. Determine *discrete logarithm* ℓ of a point Q such that

$$Q = \ell P.$$



Generic ECDLP Attacks

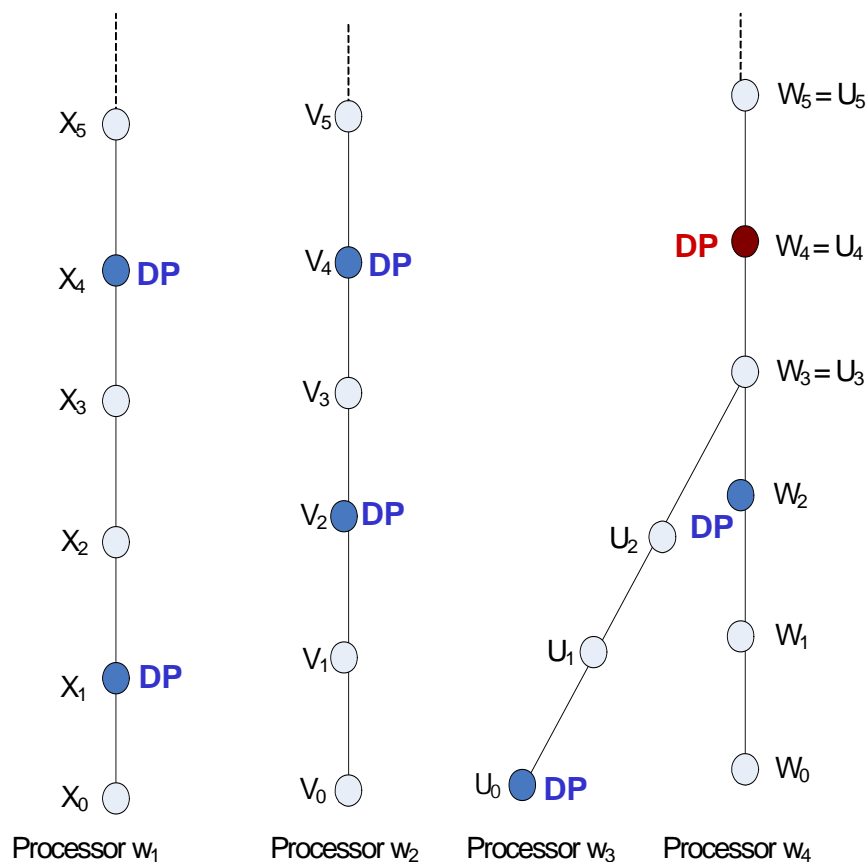
If parameters are chosen with care, only generic attacks are possible

- 1. Naïve Search:** Sequentially test $P, 2P, 3P, 4P, \dots$
 - Brute force attack is infeasible if $\#E \geq 2^{80}$
- 2. Shank's Baby-Step-Giant-Step Method**
 - Complexity in time AND memory of about $\sqrt{\#E}$
- 3. Pollard's Rho method (ρ)**
 - Most efficient algorithm for general ECDLP
 - Complexity of $\sqrt{\#E}$



Note: All attacks are *exponential* in the bit length of the group order

Multi Processor Pollard Rho (MPPR)



Colliding DP trails of multiple processors w_i

Best known attack against general ECC

Proposed by van Oorschot/Wiener in 1999

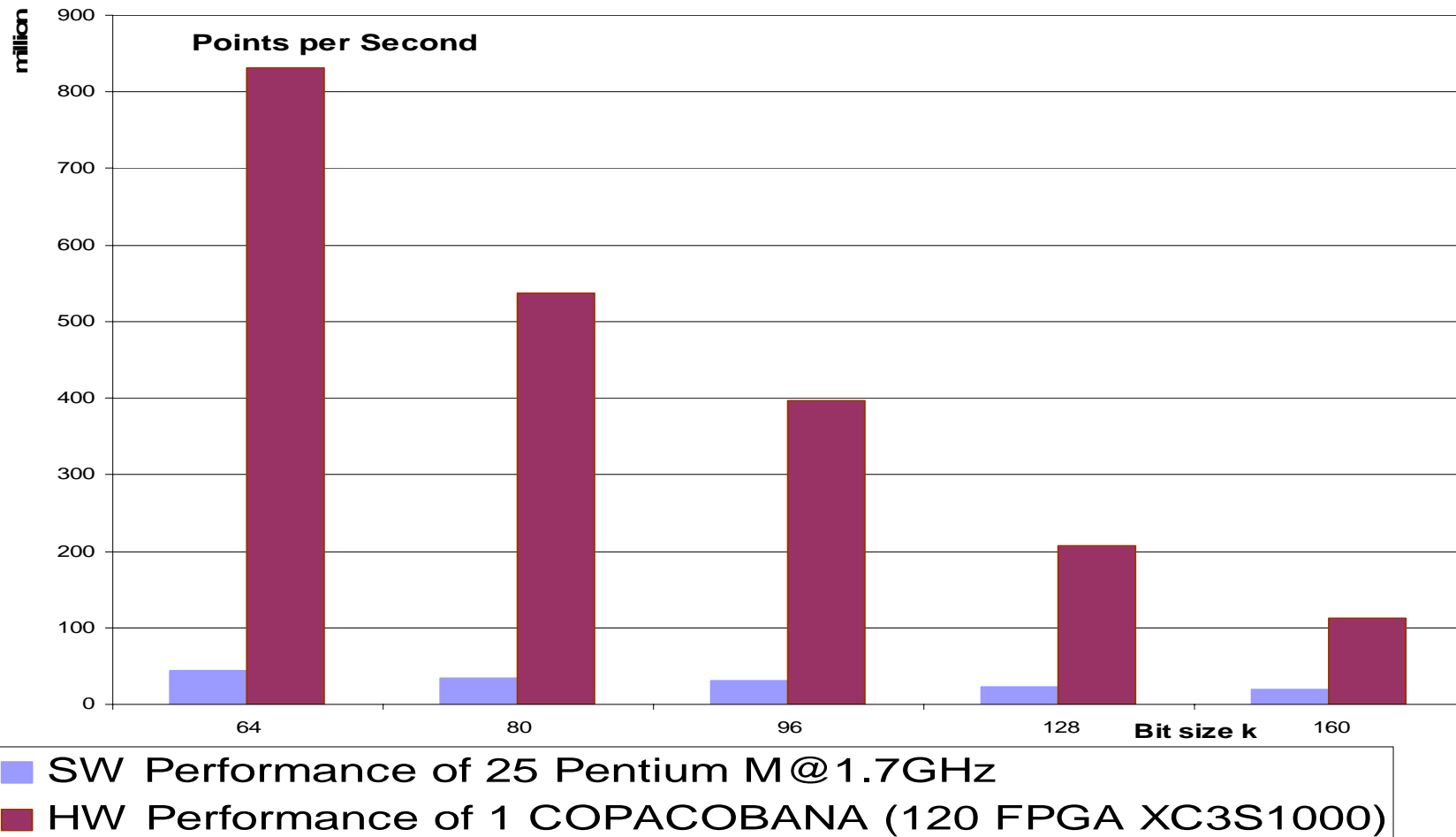
Processors have individual search paths for “Distinguished Points” (DP)

DP are stored at central server

Duplicate DP = ECDLP solution

Ideal parallelization: speed up linear in number of employed processors

ECDLP Attack Comparison: SW vs. HW for \$10.000



ECDLP Attacks for US\$ 1 million

Bit size k	SW Reference Pentium M@1.7	COPACOBANA	est. ASIC
80	40.6 h	2.58 h	-
96	8.04 d	14.8 h	-
112 (SEC-1)*	6.48 y	262 d	1.29 d
128	1.94×10^3 y	213 y	1.03 y
160	1.51×10^8 y	2.58×10^7 y	1.24×10^5 y

* SECG (STANDARDS FOR EFFICIENT CRYPTOGRAPHY)

Conclusion – COPACOBANA

- Results
 - DES in 8.6 days
 - ECCp163 attack currently \approx \$ 1 trillion ($\$10^{12}$)
 - Moore's Law: ECC 160 will stay secure for \approx 20 years
 - ECC112 (SEC-1 standard): insecure!
 - possibly real-time attack against ePass
- Many marginally weak ciphers are breakable
- „Strong“ ciphers (AES, RSA-1024, ECC-163, ...) not breakable, but robust estimates by extrapolation of COPACOBANA results
- Several future applications are currently investigated
- Pictures, papers, and much more at www.copacobana.org
- We are looking for partners for other applications