

# How to Tell the Birds from the Primes

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**Ron was wrong, Whit is right**

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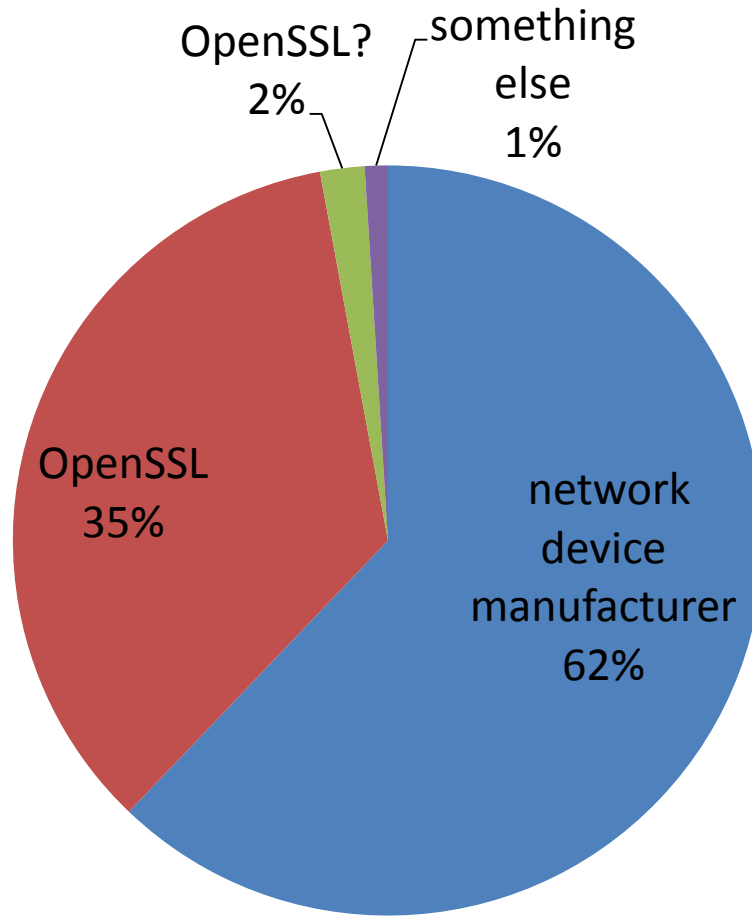
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# Ron was wrong, Whit is right, Arjen is ...?

- Snapshot of EFF Observatory from 2010
- 4.4M distinct keys
- 10K distinct factorizable RSA moduli
- Reproduced results from Lenstra et al., ePrint 2012/64

# Origin?



# How OpenSSL Samples Safe Primes?

Safe primes:  $q = 2p + 1 \in \mathbb{P}$ , where  $p \in \mathbb{P}$

## 1. Sieving step:

Find “candidate” odd  $q$ , such that  
 $q \not\equiv 0, 1 \pmod{3, 5, \dots, 17863}$

## 2. Verification step:

- a) Apply Miller-Rabin to  $q$
- b) Apply Miller-Rabin to  $(q - 1)/2$

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OpenSSL can only output a prime  $q$ ,  
such that  $q - 1 \nmid 3, 5, \dots, 17683$

1:15 primes have this property

Not a vulnerability, but a fingerprint



# Origin

