

# Physical Zero-Knowledge Proofs of Physical Properties

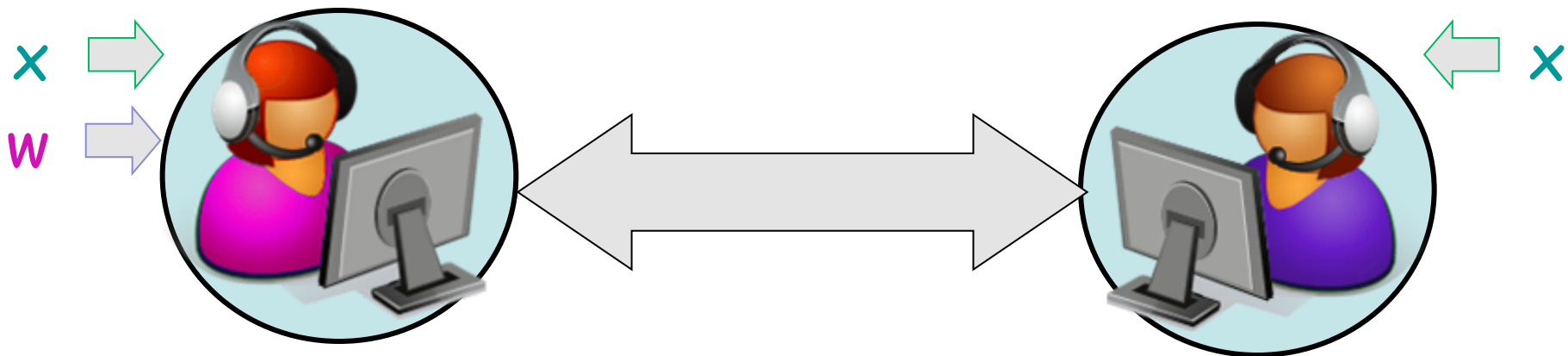
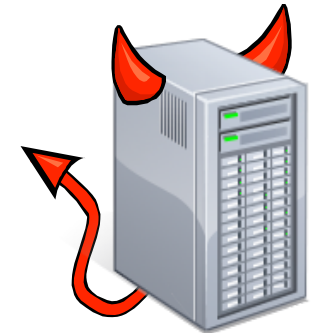
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# Typical Zero-Knowledge Scenario

- Alice and Bob receive input  $x$
- Alice has input  $w$
- Alice wants to convince Bob that
  - There is a  $w$  such that  $R(x, w) = 1$
  - She knows such a  $w$
- Alice and Bob exchange messages



# What if R is a physical property?

- Suppose the input  $x$  is physical, and R is a physical property  $\Pi$
- There is a **physical measurement** M that verifies:  
 $\Pi(x) = 1$ , i.e. “ $x$  has property  $\Pi$ ”
- Can Alice convince Bob without revealing anything more about  $x$ ?

More difficult to formalize the  
Zero-Knowledge property

# Simple Example

Alice claims she can distinguish Coke from Pepsi:



Bob selects randomly from {Coke, Pepsi}



Is it Zero-Knowledge??

What happens if Alice gets a mixture?



Which one did I give you?



Um..Coke!



If Alice **cannot** distinguish, she succeeds only with probability  $1/2$

Repeat  $t$  times

Probability Alice succeeds is  $1/2^t$

Setting not  
Inherently  
physical

## Related Work

- Physical techniques for aiding cryptographic protocols
  - Tamper-proof tokens, tamper-evident seals (envelopes), physically uncloneable functions, more examples...[GO96, GLM+04, MS08, HL08, GIS+10, GKR08, BFSK11]
- Can we find simple cryptographic protocols that humans can physically implement unaided?
  - Visual Cryptography [Naor-Shamir'94], Applied Kid Cryptography [Naor-Naor-Reingold'99], Computations with a Deck of Cards [Stiglic'01], Zero-Knowledge for Sudoku Puzzles [Gradwohl-Naor-Pinkas-Rothblum'09]
  - It's hard to see what's going on inside a computer
  - **Very relevant to voting!**
    - Polling with envelopes [Moran-Naor'06]

Inherently  
physical setting

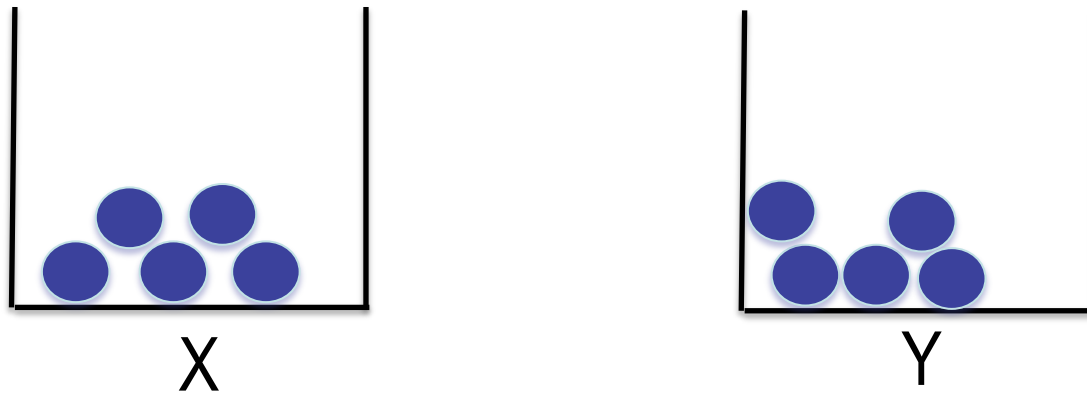
## Related Work

- Distance bounding protocols [Brands-Chaum'93]
  - Prove that you are close to a certain location
  - Use timing (speed of light)
- Boaz Barak, Alex Glaser, and Rob Goldston [GBG12] applied a zero-knowledge style technique to nuclear warhead verification
- Inherently *physical*. Not just using physical tools to construct a low complexity solution to a digital problem.

# Nuclear Warhead Verification

- Nuclear Disengagement: plan to reduce nuclear weapon stockpiles worldwide.
- START treaty, Russia and US
  - Alice promises to dismantle some of her warheads
  - How does Bob know that Alice's warhead is authentic?
  - Can Alice ensure that Bob doesn't learn, (too, much) about the design of her warhead?
- Barak et. al. reduce the problem to a protocol for Bins and Balls

# Bins and Balls



Do bins X and Y contain the same number of balls?



# This Work

- Paradigm for formally defining, modeling, analyzing physical zero-knowledge protocols
- Nuclear Disarmament: perfect physical zero-knowledge proofs for arms-control
  - Barak et. al. gave  $\epsilon$ -knowledge
- DNA Privacy: zero-knowledge proofs for DNA profiling

# Modeling physical protocols

- Separate into *logical layer* and *physical layer*
- *Physical layer*: Physical operations assumed to achieve ideal functionalities (physical assumptions)
- *Logical layer*: Hybrid world protocol obtained by replacing all physical operations with calls to their ideal functionalities.

# Modeling Example

Operation: pour  $x$  balls into a bin, and seal it

- $\mathcal{T}$  stores tuples (value, id, creator, holder, state)
- Upon receiving commands **Create**( $x$ , id) and **Seal**(id) from party  $P_i$ ,  $\mathcal{T}$  stores ( $x$ , id,  $P_i$ ,  $P_i$ , *sealed*)
- $\mathcal{T}$  only accepts **Open**(id) from the holder
- **Force**(id) causes  $\mathcal{T}$  to return entire tuple of  $id$ , and send the message “cheater” to all parties

Emulates real behavior of party that forcefully breaks open the seal without permission

# Ideal functionality $ZK^\Pi$

- Oracle access to ideal functionality  $\mathcal{M}^\Pi$
- Obtains “access” to input  $x$
- Queries  $\mathcal{M}^\Pi$  with input  $x$
- Outputs  $\Pi(x)$  to Verifier

Measurement  
verifying  $\Pi$

Full definition  
accounts for  
cheating

**Security:** Show that the *logical layer* (hybrid world translation of physical protocol) emulates  $ZK^\Pi$

# Differences from standard ZK

- No witness
  - Asymmetry between Prover and Verifier is in *access permission*, not secret knowledge or computational resources
- Ideal functionality performs verification on its own
  - It is given access permission to the input
  - Normally, Prover is required to supply a witness
- Verifier can forcefully cheat
  - Similar to covert model

# Physical ZK in UC framework

- **Benefits:**
  - Modular design and analysis of physical protocols
  - Arbitrary composition of physical and computational subprotocols
- **Feasibility:**
  - Sim does not need to do a straight-line extraction of a witness from the real world prover

# Public coin and publicly executable proofs

- *Public-coin protocols:*
  - In public-coin protocols, the verifier's messages consist only of public coin flips
  - Public-coin physical protocols are **publicly executable**
  - The verifier can sit behind a glass screen throughout the execution

**Makes a huge difference  
for physical security!**

- **We construct a publicly executable DNA inequality protocol**

# Feasibility of publicly executable proofs?

- In the standard digital setting, public-coin ZK = private-coin ZK [Oka96, GSV98, GV99, Vad04]
- General result for physical zero-knowledge?
- Techniques for explicit conversions of private-coin protocols to public-coin protocols don't translate well in the physical setting
  - Universal hashing of physical messages?
  - Physically concealed messages

**Public coin => publicly executable**  
**Publicly executable =>? public coin**



# Summary and Further Research

Physical Cryptography is

- Relevant
- Fun
- **Structured(?)**: connections with known crypto/complexity techniques
- Many foundational questions remain