

# Analysis of Step-Reduced SHA-256

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# SHA-256 is Interesting and Challenging

FIPS Standard since 2002

Option for a SHA-1 upgrade

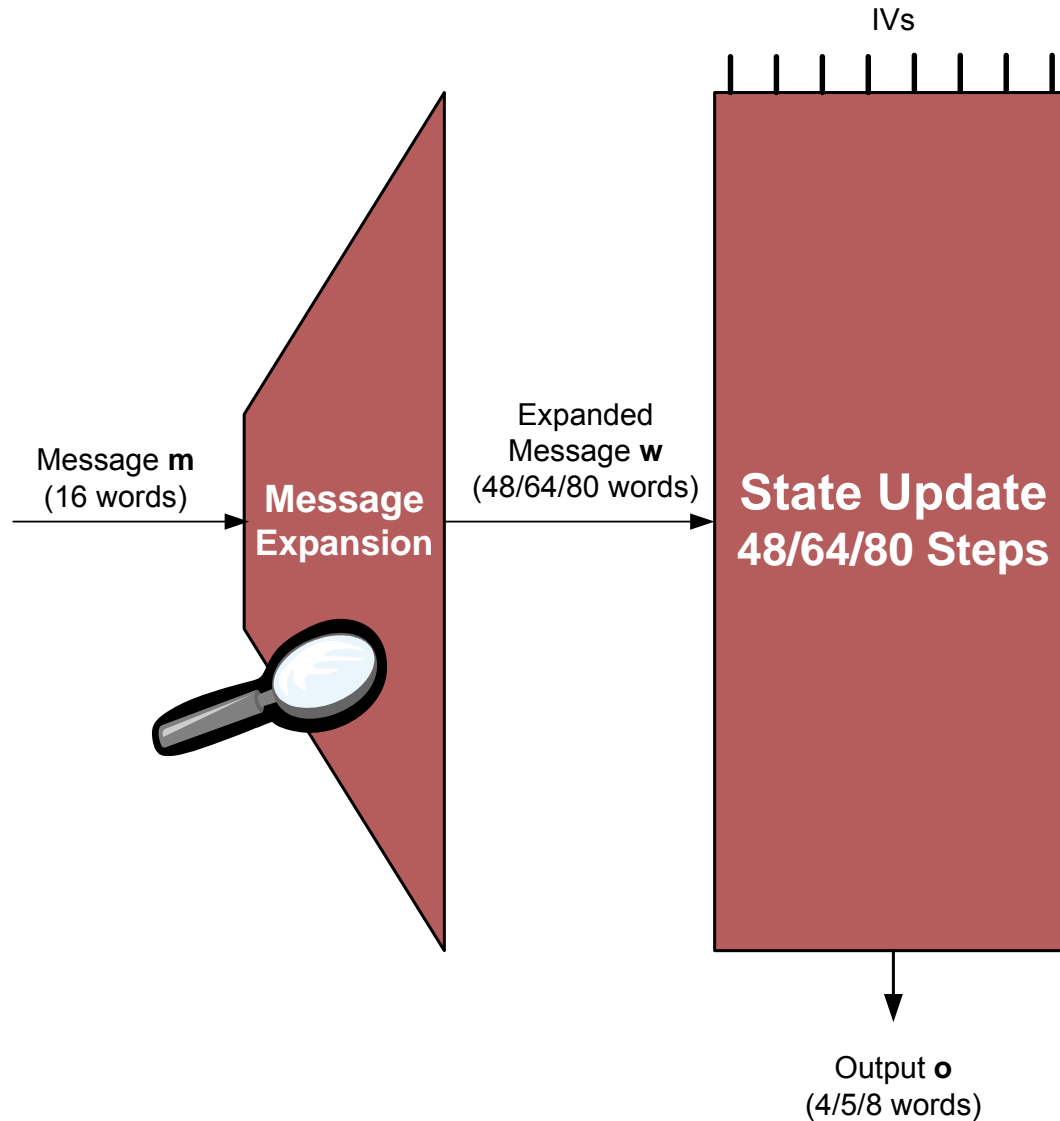


Prudent to know:

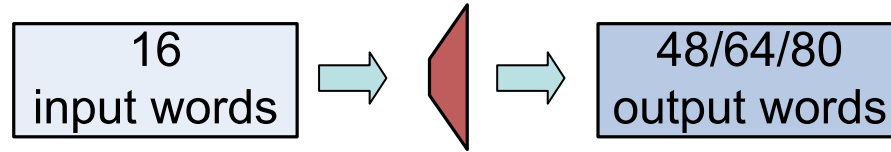
How hard is it to find collisions for SHA-256?

What about step-reduced variants (security margin)?

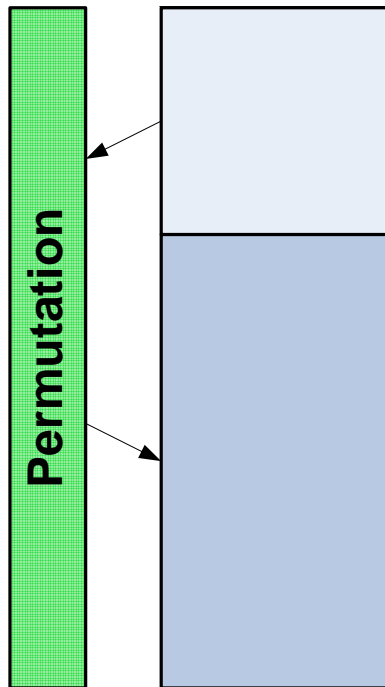
# Outline of MD4-style Hash Functions



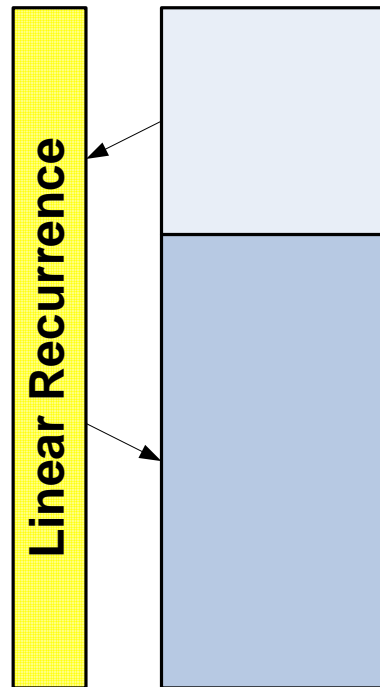
# Message Expansions in the MD4 family



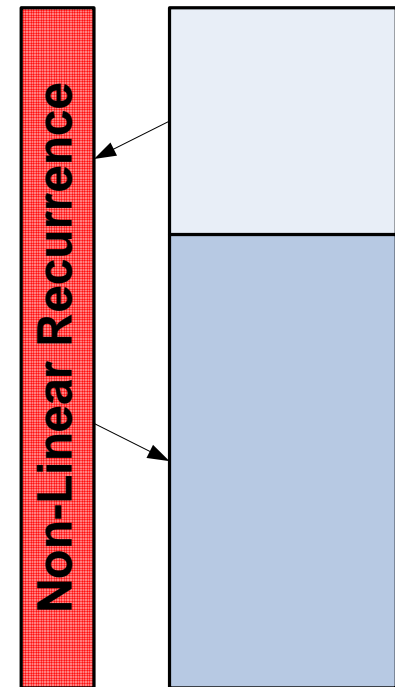
**MD4/5, RIPEMD**



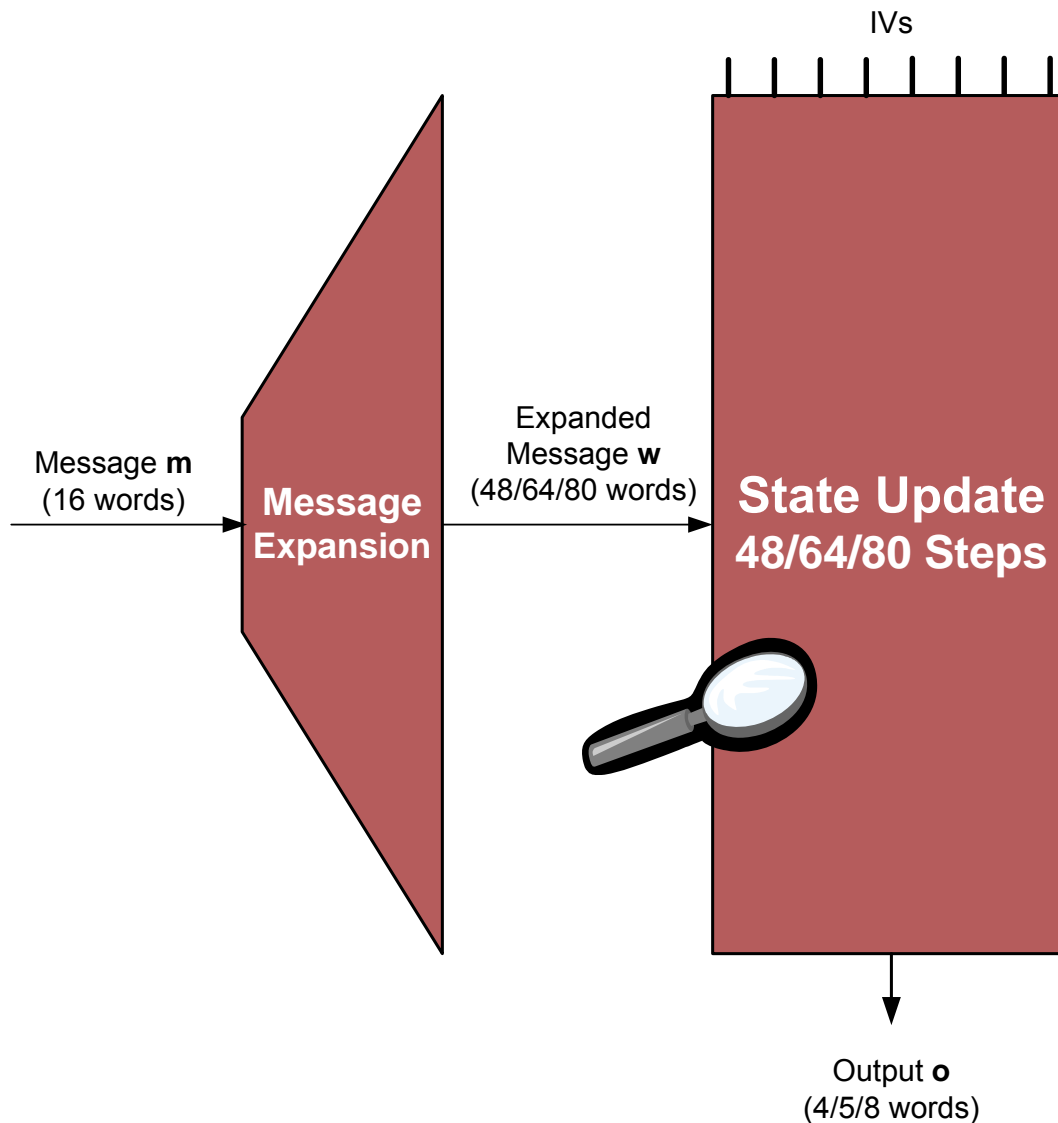
**SHA-0 / SHA-1**



**SHA-2 family**

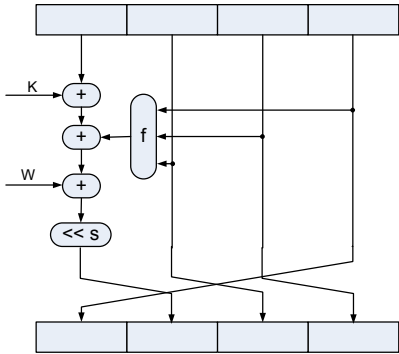


# Outline of MD4-style Hash Functions

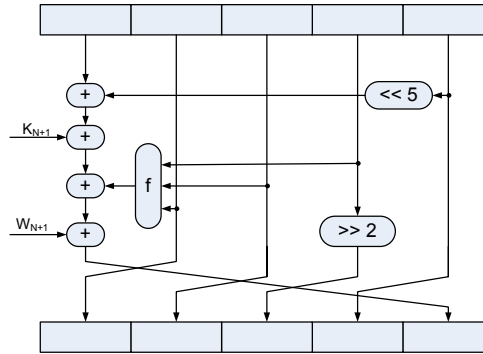


# Evolution of the State Updates in the MD4 Family

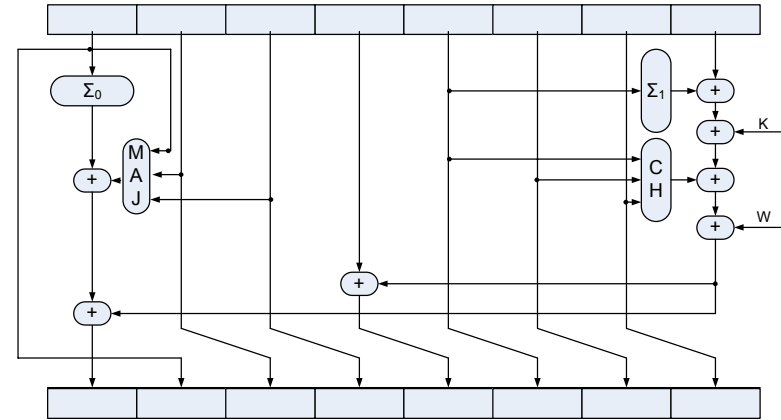
MD4



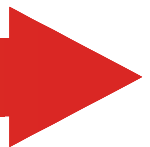
SHA-0/1



SHA-2 family



Design Complexity



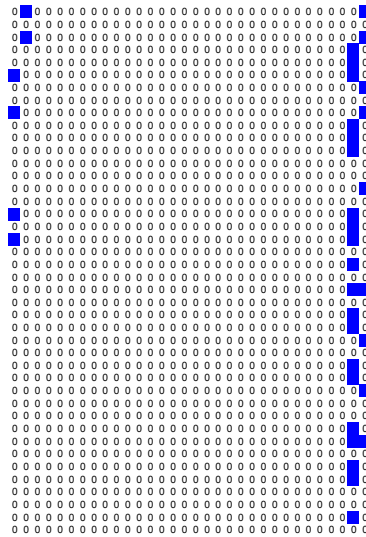
# Overview

- Top-level review of results on SHA-1
- Applicability to SHA-2 members
- New method overcomes identified obstacles
- Interesting insights and directions for future work

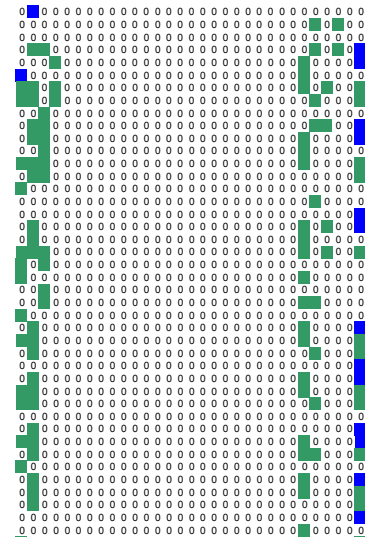
# Review of Collision Attacks on SHA-1

[CJ98, BC04, RO05, BCJ05, WYY05]

step 1



apply corrections

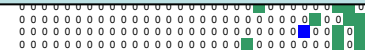


- perturbation
- correction

Two properties are needed for that:

- ME is invariant with respect to rotation
- ME is invariant with respect to translation

step 80



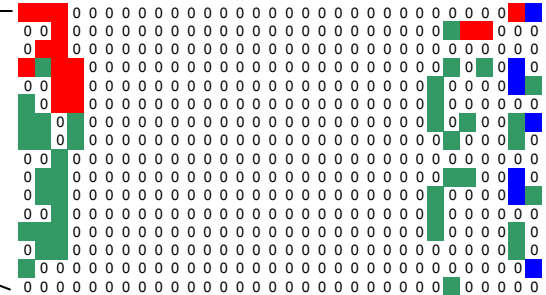
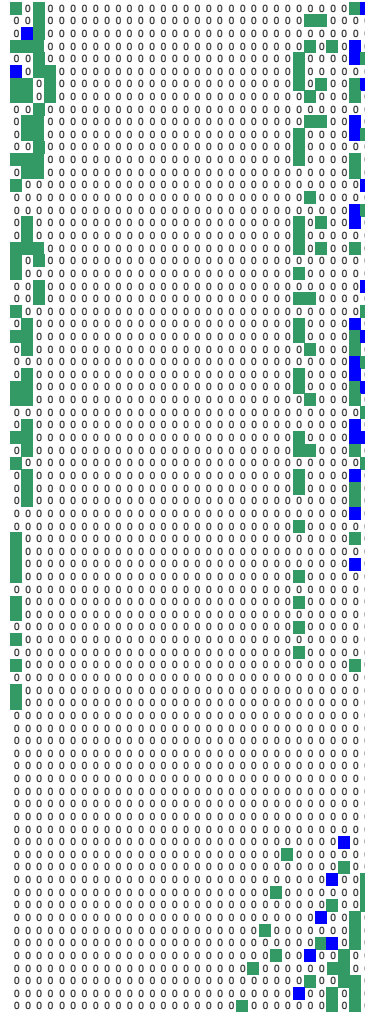
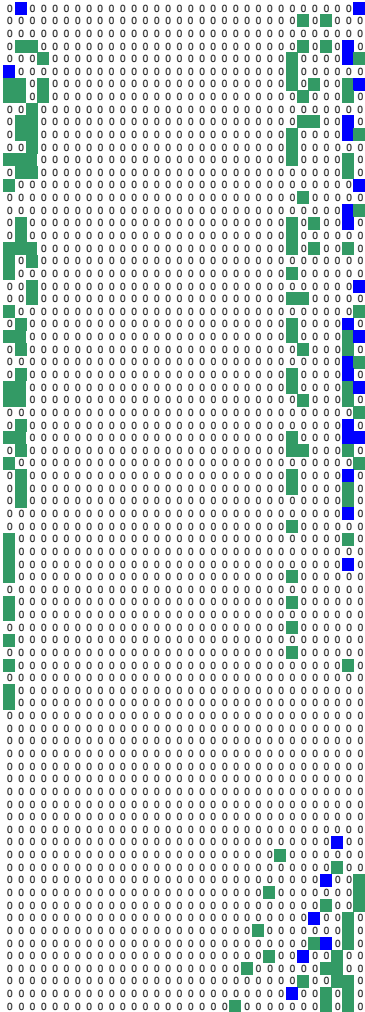





# needed

# reality

# difference

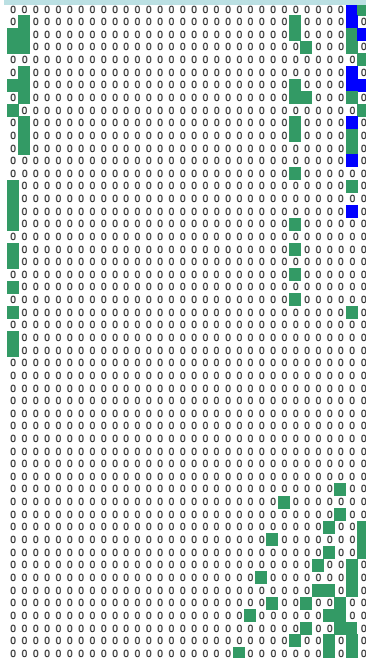


 *ghost differences of type 1*

# Review of the [WYY05] Characteristic

Message  
Modification  
improves the  
probability to  
 $2^{-6x}$

**Low-probability** characteristic

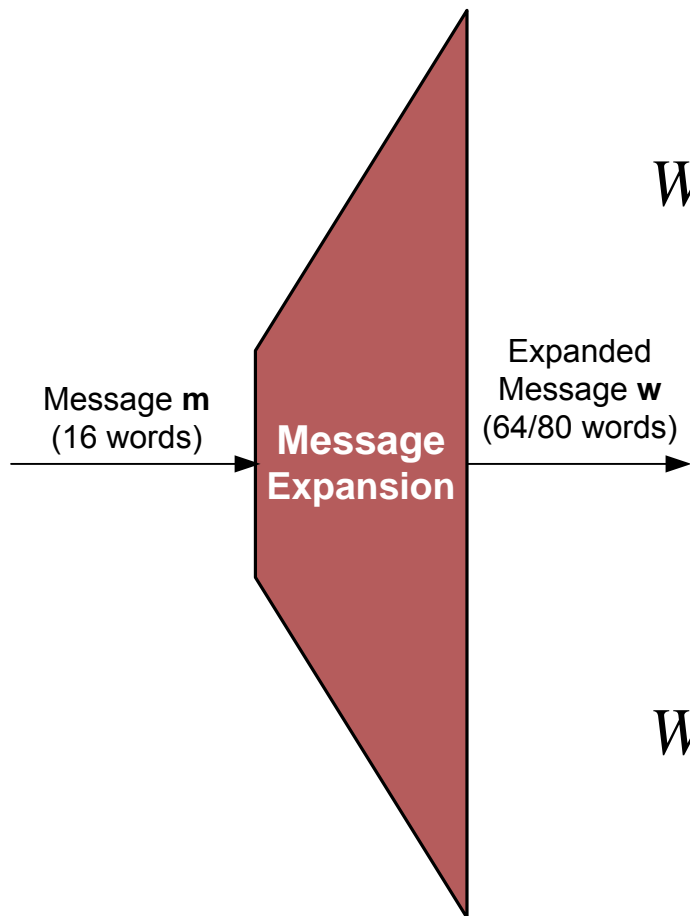


**High-probability** characteristic (about  $2^{-83}$ )

# Comparison of SHA Message Expansions

## SHA-1

$$W_t = \begin{cases} M_t & \text{for } (0 \leq t \leq 15) \\ \text{ROTL}^1(W_{t-3} \oplus W_{t-8} \oplus W_{t-14} \oplus W_{t-16}) & \text{for } (16 \leq t \leq 79) \end{cases}$$



## SHA-256

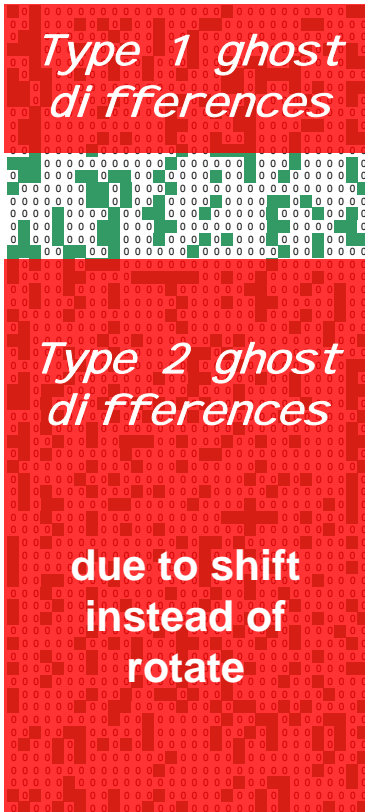
$$W_t = \begin{cases} M_t & \text{for } (0 \leq t \leq 15) \\ \sigma_1(W_{t-2}) + W_{t-7} + \sigma_0(W_{t-15}) + W_{t-16} & \text{for } (16 \leq t \leq 63) \end{cases}$$

$$\sigma_0(x) = \text{ROTR}^7(x) \oplus \text{ROTR}^{18}(x) \oplus \text{SHR}^3(x)$$

$$\sigma_1(x) = \text{ROTR}^{17}(x) \oplus \text{ROTR}^{19}(x) \oplus \text{SHR}^{10}(x)$$



# Approach does not apply to SHA-2



Low-probability characteristic

~~High-probability characteristic~~  
Low-probability characteristic

# Carry Effects in the Message Expansion

What about the non-linearity of the Message Expansion?

Theorem:  
Preventing *type 2 ghost differences* is  
not always possible

## Approach to avoid *type-2 Ghost Differences*

- Build up on approach originally pioneered by Rijmen and Oswald [RO05]
- Generalization  $\rightarrow$  huge search space  $2^{768}$
- **Solution:**
  - Generic + heuristic search-space reduction  $\rightarrow 2^{64}$
  - Probabilistic search using algorithms from coding theory [Leo88,CC98]

# Example of 19-step Characteristic

Step	W'	A'	B'	C'	D'	E'	F'	G'	H'
1-4	0	0	0	0	0	0	0	0	0
05	85009008	85009008	0	0	0	85009008	0	0	0
06	a14cae12	a1442610	85009008	0	0	02000802	85009008	0	0
07	0	0	a1442610	85009008	0	084c4120	02000802	85009008	0
08	8200a8a8	00000020	0	a1442610	85009008	00000020	084c4120	02000802	85009008
09	85009008	85009008	00000020	0	a1442610	01008008	00000020	084c4120	02000802
10	0	0	85009008	00000020	0	02000802	01008008	00000020	084c4120
11	0	0	0	85009008	00000020	0	02000802	01008008	00000020
12	0	00000020	0	0	85009008	0	0	02000802	01008008
13	0	0	00000020	0	0	84001000	0	0	02000802
14	00088802	0	0	00000020	0	0	84001000	0	0
15	0	0	0	0	00000020	0	0	84001000	0
16	0	0	0	0	0	00000020	0	0	84001000
17	0	0	0	0	0	0	00000020	0	0
18	0	0	0	0	0	0	0	00000020	0
19	0	0	0	0	0	0	0	0	00000020

**1-block collision for SHA-224**

## Interesting Results

- **Perturbation pattern** is **no valid expanded message**
  - But the sum of perturbations and corrections is
  
- **More freedom** for the **carry**
  - ... to prevent impossible characteristics
  
- The **overall probability** is **much higher** than the product of the probabilities of each individual local collision
  - Different to SHA-0 / SHA\_1
  - Example: low-weight 19-step characteristic
    - 23 local collisions of probability around  $2^{-40}$
    - Total probability is much higher: instead of  $2^{-920}$  around  **$2^{-200}$**   
 (Compare this to a similar probability of the best known 80-step characteristic for SHA-1)



# Conclusions

- First analysis of unmodified SHA-256/224 for a nontrivial number of steps
- Collision resistance of SHA-256/224 is not threatened
- All publicly known attacks on SHA-0/1 since 1997 are not directly applicable to any SHA-2 member
- New analysis method
  - Circumvent problem of *ghost differences of type 2*
  - New type of perturbation pattern
  - Probability of a local collision is much less relevant
  - Explicit control of carry extensions is possible and needed

# Future Research

1. Ways to reduce the search space for high probability characteristics
2. New message modification techniques
3. Exploiting non-linearity of Message Expansion
4. Apply multi-block approach

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