Complete Cover deCryption: the Challenge of LNCS 6805

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Rump session EUROCRYPT '12 Cambridge, UK

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1233



Festschrift

LNCS 6805

Cryptography and Security: From Theory to Applications

Essays Dedicated to Jean-Jacques Quisquater on the Occasion of His 65th Birthday

David Naccache (Ed.)







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Festschrift

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1. David Naccache (Ed.)

- He did the job finding and editing
- 1 + 32 papers
- With 84 authors (12 in the room)

Naccache (Ed.)

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Festschrift

Cryptography and Security: From Theory to Applications

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2. LNCS Festschrifts

- Festschrifts honor individual researchers and their scientific work, or they honor institutions or fields.
- Historical and even personal aspects may show up.
- They present internationally relevant technical contributions with a reasonable topical focus.
- Designed with a fresh orange and blue cover.



3. LNCS 6805

• This number is also the number of a processor family (Motorola) used for the first smart cards I was working

3. LNCS 6805

Motorola - 6805		
6805 Family		
Homepage: none ?		
Family description:		
Designer: Motorola	Architecture: 8-bit	
Initial release: ?	Frequencies: ?	
Datasheet: ?	Technology: ?	

Motorola - MC6805R2L1

MC680SR2L1

General Specifications:					
Manufacture:	Motorola	Family:	6805		
Chip type:	MCU	Introduced:			
Speed:		Architecture:	8-bit		
Application:					
Architecture Specifications:					
CPU arch:		ISA:			
Microarch:					
Processor core:		# of cores:			
Designer:	Motorola	FPU:	NA		
Ext data bus:	8	Address bus:	8		
Instruction Set:					
Features:					
Technology:					
Technology:	HMOS	Process:			
Transistors:		Die size:			
Vcc:		Voltage I/O:			

• This number is also the number of a processor family (Motorola) used for the first smart cards I was working

6800 or 6805

- David reserved long time ago 2 numbers from LNCS: 6800 and 6805
- I did the choice
- Waiting 8051(the main family of processors I was working) was too long ...

My only contribution on June 3, 2011





COprocessor for RSA In a Rush

CORSAIR: A Smart Card for Public Key Cryptosystems

Dominique de Waleffe & Jean-Jacques Quisquater

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Abstract. Algorithms best suited for flexible smart card applications are based on public key cryptosystems — RSA, zero-knowledge protocols ... Their practical implementation (execution in ≈ 1 second) entails a computing power beyond the reach of classical smart cards, since large integers (512 bits) have to be manipulated in complex ways (exponentiation). CORSAIR achieves up to 40 (8 bit) MIPS with a clock speed of 6 Mhz. This allows to compute $X^E \mod M$, with 512 bit operands, in less than 1.5 second (0.4 sec for a signature). The new smart card is in the final design stage; the first test chips should be available by the end of 1990.

Keywords: smart card, public key algorithms, RSA, digital signature, zero-knowledge protocols.

1 Introduction

A large number of security problems can be solved by correct use of cryptographic methods. However, all methods found to date are more or less computationally intensive.

- DES works by applying a complex multiround algorithm on medium size numbers.
- Diffie-Hellman key exchange protocol is based on modular exponentiation of large integers.
- RSA is based on the same exponentiation and needs large exponents.
- Zero-knowledge protocols like those of Fiat-Shamir [9] or Guillou-Quisquater [10] use large number exponentiation but the exponents are not as large as in RSA.
- Many identity-based systems also rely on modular exponentiation of large numbers.

Public key techniques are the most promising for the future as they provide more flexible solutions and impose less burden both on users and security management. Most practical techniques rely on large integer arithmetic.

A.J. Menezes and S.A. Vanstone (Eds.): Advances in Cryptology - CRYPTO '90, LNCS 537, pp. 502-513, 1991.
© Springer-Verlag Berlin Heidelberg 1991

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Figure 2: Simple cell



Figure 5: Final architecture

Next step: FAME from NXP-Philips



Using photoshop (fuzzy) and adding









Secret key



Secret key





Secret key









How to Explain Zero-Knowledge Protocols to Your Children

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in collaboration with Tom BERSON(3) for the English version

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The Strange Cave of Ali Buba

◇ Know, oh my children, that very long ago, in the Eastern city of Baghdad, there lived an old man named Ali Baba. Every day Ali Baba would go to the bazaar to buy or sell things. This is a story which is partly about Ali Baba, and partly also about a cave, a strange cave whose secret and wonder exist to this day. But I get ahead of myself ...

One day in the Baghdad bazaar a thief grabbed a purse from Ali Baba who right away started to run after him. The thief fled into a cave whose entryway forked into two dark winding passages: one to the left and the other to the right (The Entry of the Cave).

Ali Baba did not see which passage the thief ran into. Ali Baba had to choose which way to go, and he decided to go to the left. The left-hand passage ended in a dead end. Ali Baba searched all the way from the fork to the dead end, but he did not find the thief. Ali Baba said to himself that the thief was perhaps in the other passage. So he searched the right-hand passage, which also came to a dead end. But again he did not find the thief.



"This cave is pretty strange," said Ali Baba to himself, "Where has my thief gone?"

The following day another thief grabbed Ali Baba's basket and fled, as the first thief had fled, into the strange cave. Ali Baba pursued him, and again did not see which way the thief went. This time Ali Baba decided to search to the right. He went all the way to the end of the right-hand passage, but he did not find the thief. He said to himself that, like the first thief, the second thief had also been lucky in taking the passage Ali Baba did not choose to search. This had undoubtedly let the thief leave again and to blend quietly into the crowded bazaar.

The days went by, and every day brought its thief. Ali Baba always ran after the thief, but he never caught any of them. On the fortieth day a fortieth thief grabbed Ali Baba's turban and fled, as thirty-nine thieves had done before him, into the strange cave. Ali Baba yet again did not see which way the thief went. This time Ali Baba decided to search the left-hand passage, but again he did not find the thief at the end of the passage. Ali Baba was very puzzled.

He could have said to himself, as he had done before, that the fortieth thief had been as lucky as each of the other thirty-nine thieves. But this explanation was so



- J.-J. Quisquater et D. Samyde:
- « a new tool for non intrusive analysis of smart cards based on electro-magnetic emissions, the SEMA and DEMA methods »
- Presented at the rump session of EUROCRYPT '2000, Bruges, Belgium.

Eurocrypt 2000

Bruges (Brugge), Belgium, May 14-18, 2000

Eurocrypt 2000 Rump Session

The occasional drink and poster session (part one)

Efficient Protocols from Homomorphic Threshold Cryptography	Ivan Damgård, Ronald Cramer, Jesper Buus Nielsen, Mads Jurik
Elliptic Curve Systems Too Risky? Or TRoublesome?	Arjen K. Lenstra
The Schoof-Elkies-Atkin algorithm in characteristic 2 - The Previous world record	Frederik Vercauteren
A New Record in point counting on elliptic curves	Pierrick Gaudry
A new tool for non-intrusive analysis of smart cards based on electro-magnetic emissions. The SEMA and DEMA methods	Jean-Jacques Quisquater, David Samyde
On the Soundness of Girault's Scheme	Fabrice Boudot
The NESSIE Call for Cryptographic Algorithms	Eli Biham
FPGA Implementation of Modular Exponentiation Using Montgomery Method	Elena Trichina
One-round secure computation and secure Autonomous Mobile Agents	Christian Cachin, Jan Camenisch, Joe Kilian, Joy Müller
	The difference of the statistic term and the second of the statistic second of the statistic second of the second of the

The occasional drink and poster session (part two)

전에서 가운 가 다시가 다 한 사람이 가 다 가지, 가운 가 다시가 다 한 가 다시가 다 한 가 다 가 다운 가 다시가 다 가 다 가 다 가 다 가 다 가 다 가 다 가 다 가 다	가 아버지가 만들어졌다. 김가 아버지가 만들어 있는 것을 가지 않는 것을 가지 않는 것을 가지 못했다. 감독
Braid Group Cryptosystem, the Arithmetic Key Agreement Protocol	Jim Hughes
Update on UMAC Fast Message Authentication	Phil Rogaway
Small generic hardcore subsets for the discrete logarithm: short secret DL-keys	Clauss P. Schnorr
A popular protocol whose security decreases as key size increases	David Naccache
Necessary and Sufficient Assumptions for Non-Interactive Zero-Knowledge Proofs of Knowledge for all NP relations	Alfredo De Santis, Giovanni Di Crescenzo, Giuseppe Persiano
A proven secure tracing algorithm for the optimal KD traitor tracing Scheme	Kaoru Kurosawa, Mike Burmester, Yvo Desmedt
Efficient Algorithms for Differential Probability modulo 2^n and Related Problems	Helger Lipmaa, Shiho Moriai

Eurocrypt 2000 Poster Session

OCIATION FOA



Inca khipus



Inca khipus





Fig. 3 - Formas de hacer nudos de un quipu, según Locke 1978 [1923]: a) Cuerda principal; b) Lazo para atar las cuerdas colgantes a la principal; c) Lazo ajustado; d) Cuerda subsidiaria, con resumen de cuentas; e) Nudo largo (no más de 9 lazos), sin ajustar; f) Nudo ajustado; g) Nudo sin ajustar; h) Nudo ajustado; i) Nudo sin ajustar; j) Nudo ajustado; k) Nudo del extremo inferior de la cuerda, sin significado numérico; l) Nudo en cuerda colgante de una cuerda subsidiaria; 1 a 9: numerales representados en los nudos.

Inca khipus













Again!



















Thanks, David!



A success!



Now the challenge!

• What is the integer value of the secret key?

• 4 digits: some hint was already given

 Give me a piece of paper with the value and your name: the prize will be given to the first correct answer or the closest one (my definition)

Now the prize!



After the rump session News from the challenge

- The winner is Mark Manulis (University of Surrey)
- My hope was that somebody remarks an anomaly on the first slide ... ☺
- The secret key is: 1233
 - See the first slide (with a fake page number) and the key:



Next game

• At CRYPTO 2012?