# On the Soundness of Girault's Scheme 

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The Girault protocol [Eurocrypt'91] is a zero-knowledge like scheme which allows Alice to prove to Bob that she knows $x$, a discrete logarithm of $y$ in base $g$ modulo $n$, where $n$ is a composite number whose factorization is unknown. It runs as follows: Alice randomly selects $r \in[0, A]$ (where $A \gg n)$ and sends $W=g^{r} \bmod n$ to Bob. Bob sends to Alice a random challenge $c \in[0, k]$. Alice sends in reply $D=r+x c(i n Z)$. Bob accepts the proof if $W=g^{D} y^{-c} \bmod n$.

Poupard and Stern [Eurocrypt'98] prove that this protocol is sound if computing discrete logarithms modulo $n$ is hard, but only when this protocol is used for identification or signature schemes, i.e. when an attacker cannot choose the public data $y$. For other contexts (when an attacker can choose the public data $y$ ), Fujisaki and Okamoto [Crypto'97] "prove" that this protocol is sound if the strong RSA assumption holds, i.e. an attacker cannot succeeds if he does not know a discrete logarithm of $y$.

We show that Fujisaki-Okamoto's claim is wrong, and hence so is their proof. Moreover, an attacker can succeed to Girault's protocol without knowing the discrete logarithm of $y$ : he can succeed with probability $1 / 2$ if he only knows the discrete logarithm of $-y$ in base $g$ modulo $n^{1}$ : if $c$ is even, the attacker replies $D=r+x^{\prime} c$, where $x^{\prime}$ is the discrete logarithm of $-y$.

To design a sound protocol, we slightly transform Girault's protocol into a zero-knowledge like proof of knowledge of a discrete logarithm of $\pm y$ in base $g$ modulo $n$. Moreover, we prove that this protocol is sound if the RSA assumption holds (instead of the strong RSA assumption).

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[^0]:    ${ }^{1}$ Note that when the factorization of $n$ is unknown, the knowledge of a discrete logarithm of $-y$ does not allow to compute a discrete logarithm of $y$ (it is not the case if $n$ is prime or when the factorization of $n$ is known). Moreover, the knowledge of discrete logarithms of $y$ and $-y$ allows to factor $n$.

