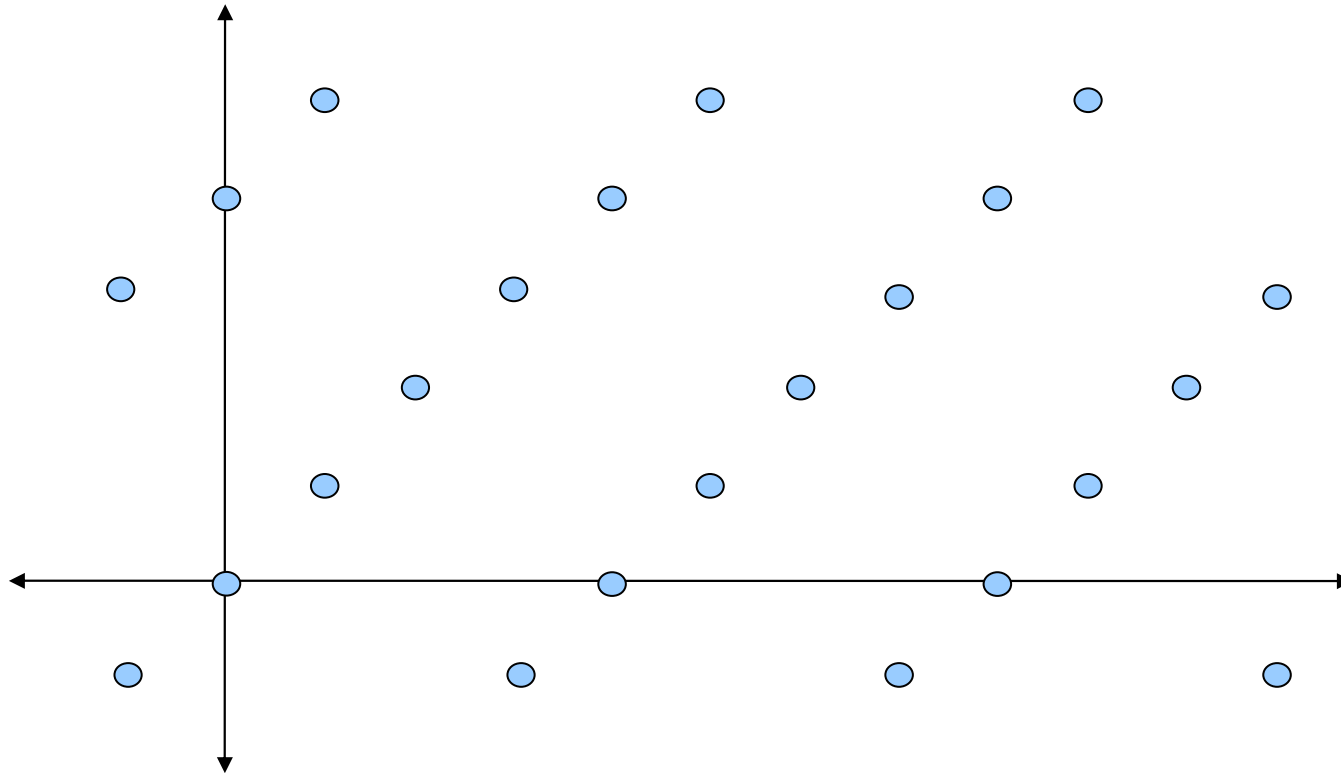


# On Bounded Distance Decoding, Unique Shortest Vectors, and the Minimum Distance Problem

Vadim Lyubashevsky

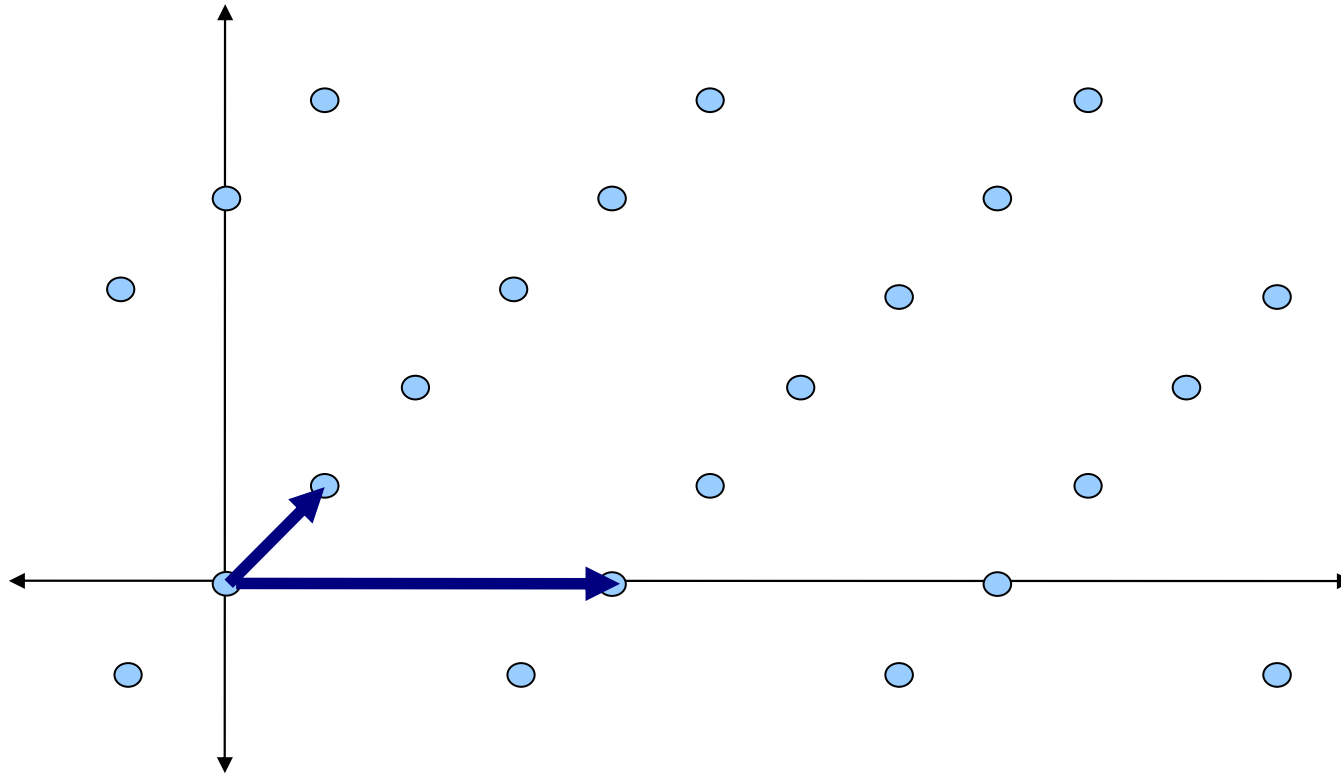
Daniele Micciancio

# Lattices



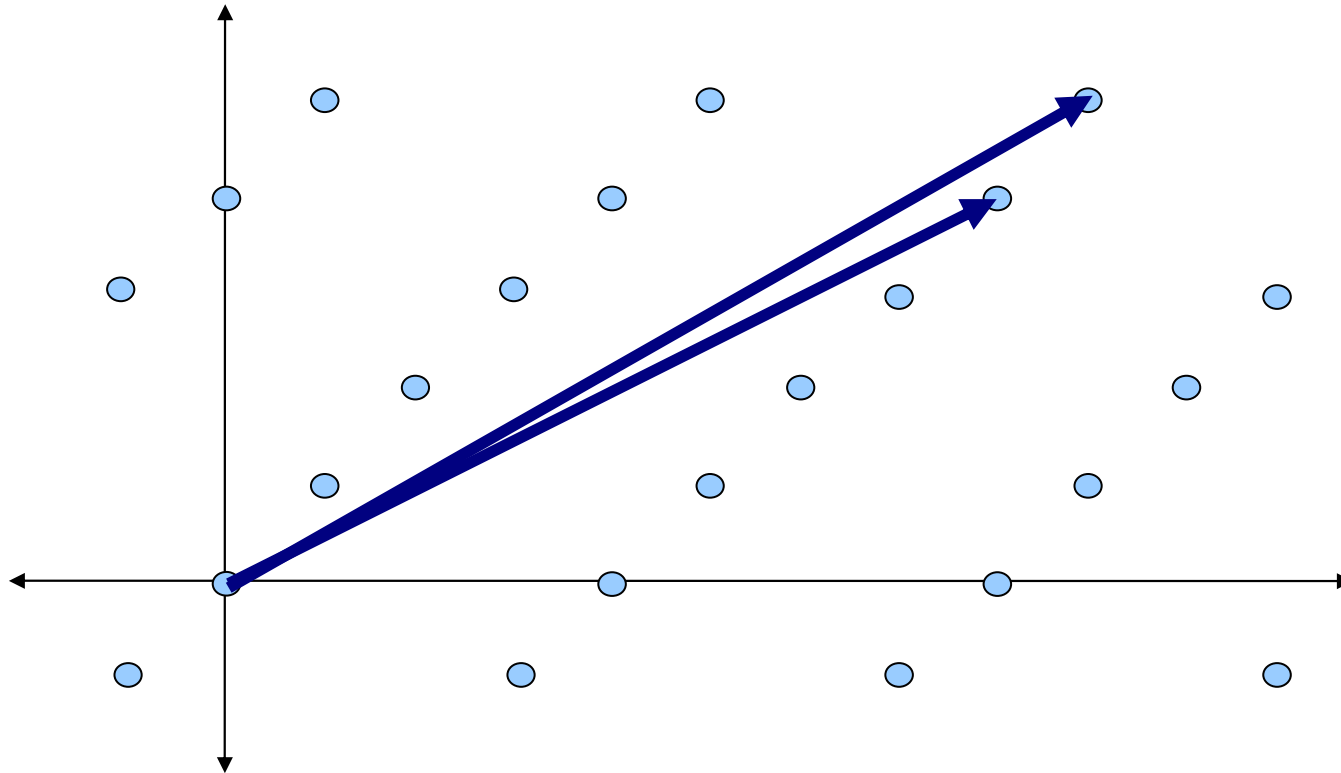
Lattice: A discrete additive subgroup of  $\mathbb{R}^n$

# Lattices



Basis: A set of linearly independent vectors that generate the lattice.

# Lattices

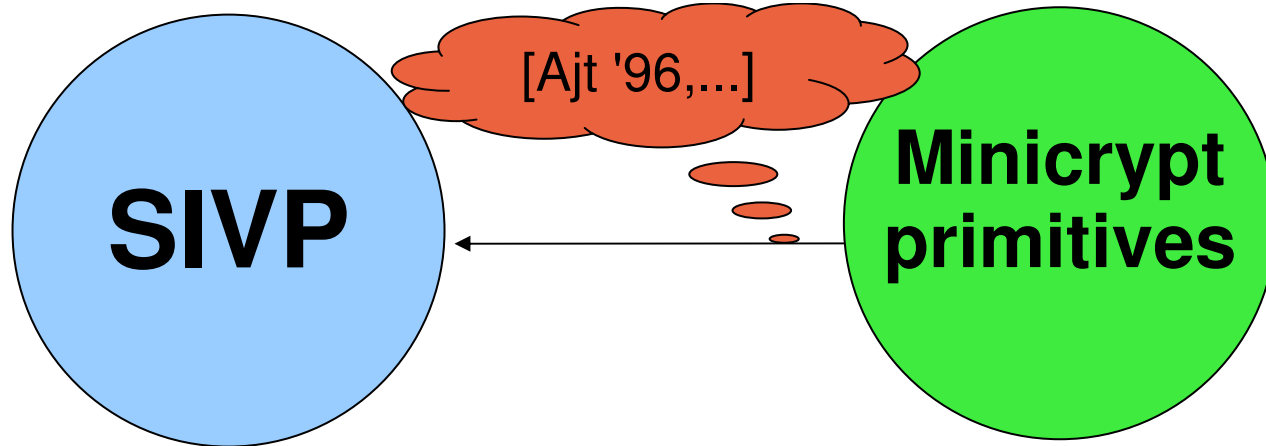


Basis: A set of linearly independent vectors that generate the lattice.

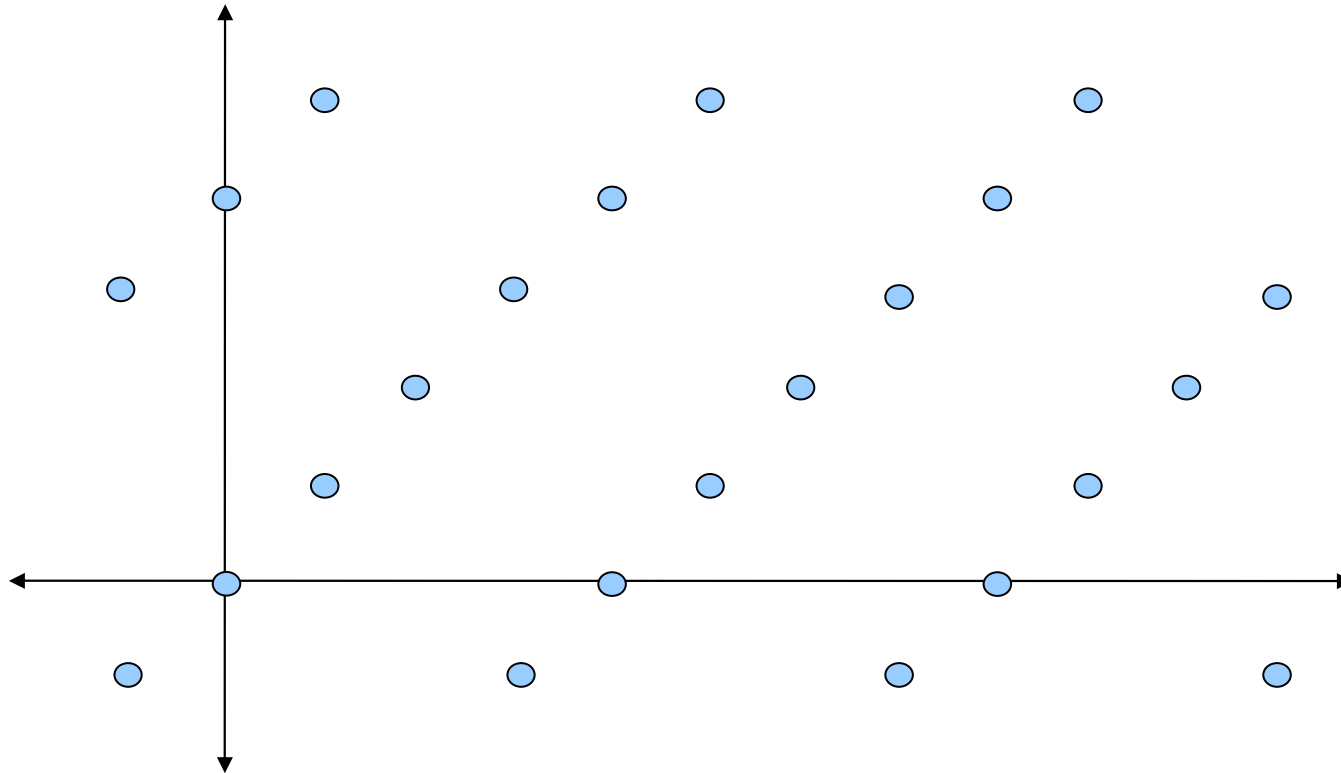
# Why are Lattices Interesting?

(In Cryptography)

- Ajtai ('96) showed that solving “*average*” instances of some lattice problem implies solving *all instances* of a lattice problem
- Possible to base cryptography on worst-case instances of lattice problems

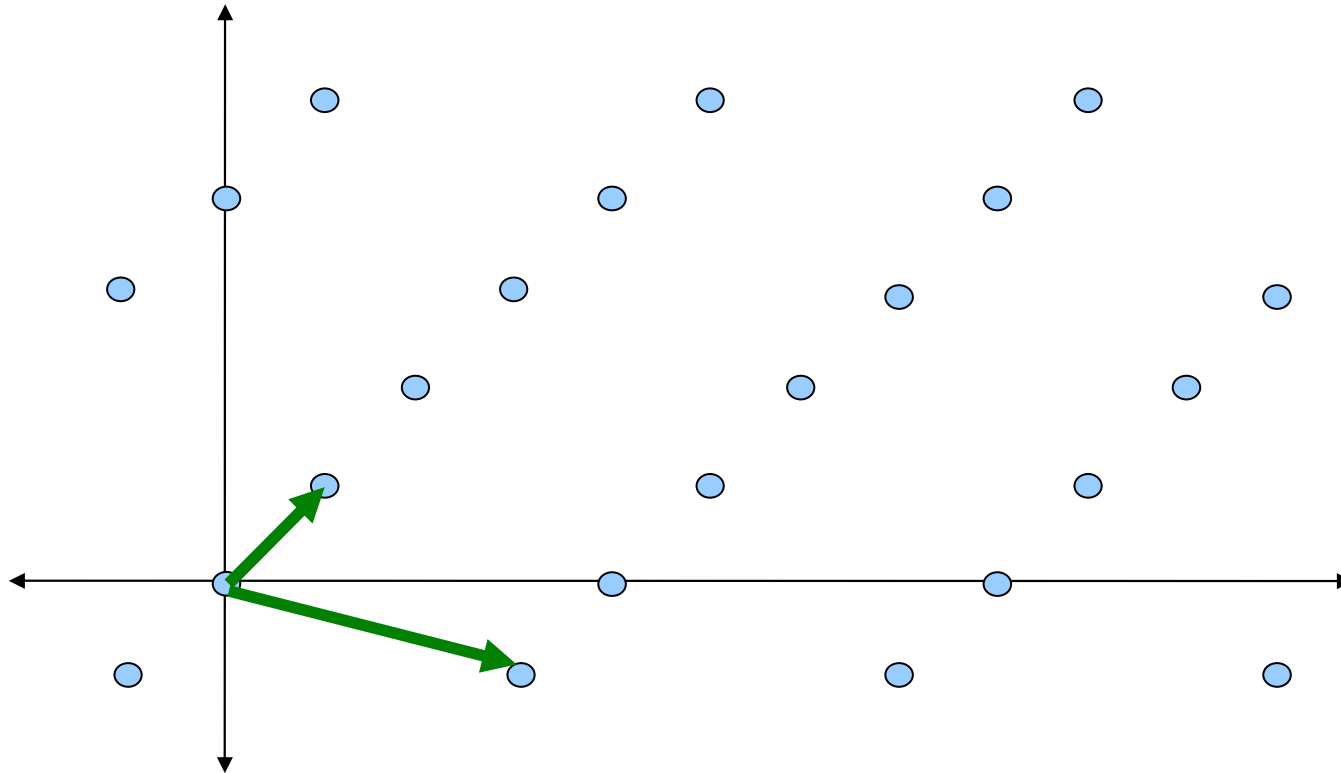


# Shortest Independent Vector Problem (SIVP)



Find  $n$  short linearly independent vectors

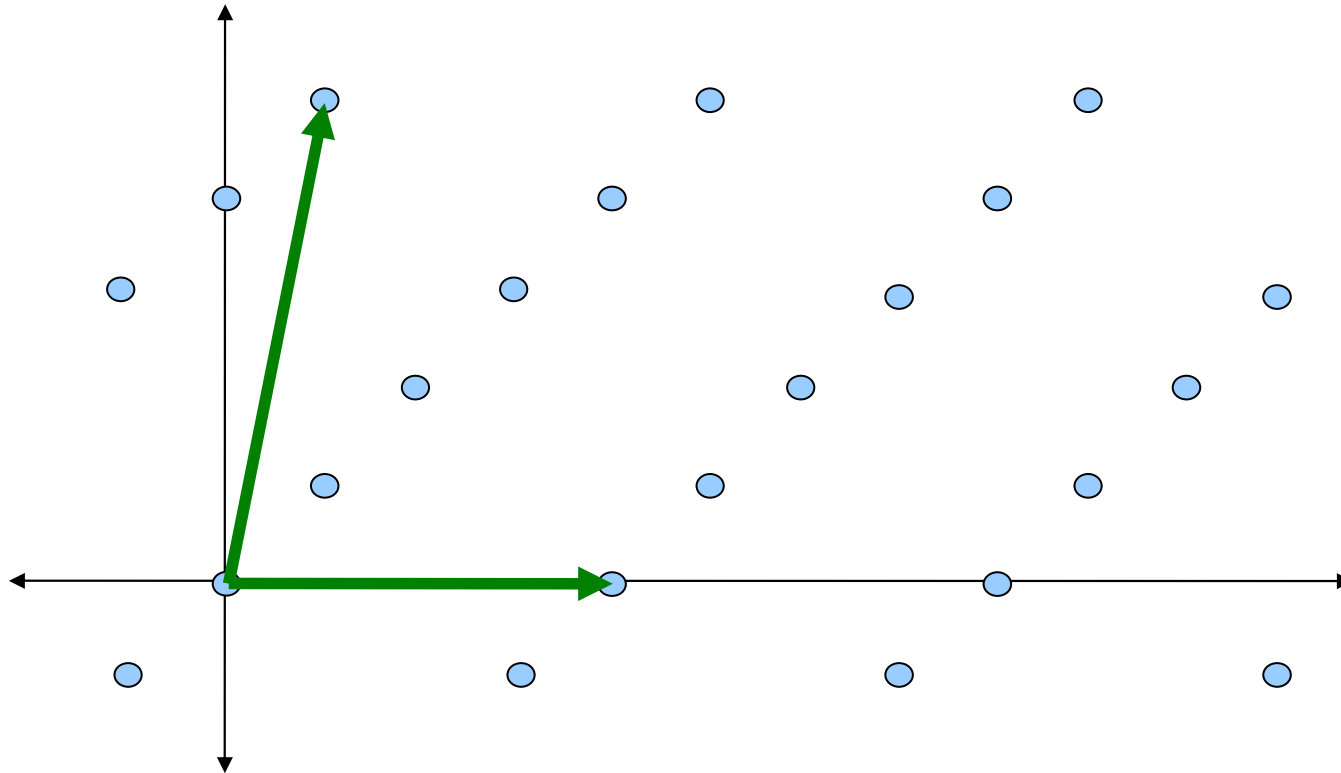
# Shortest Independent Vector Problem (SIVP)



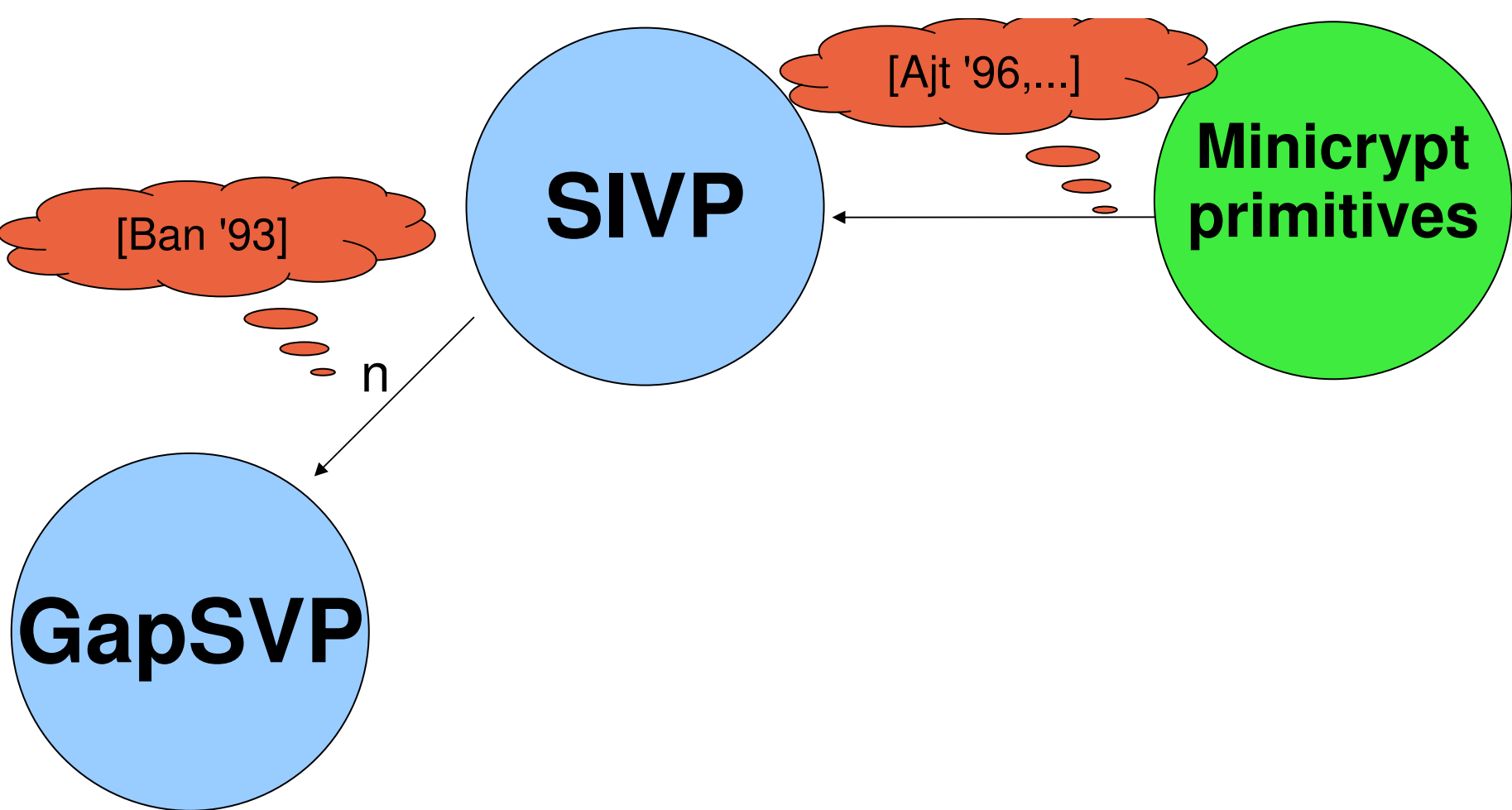
Find  $n$  short linearly independent vectors



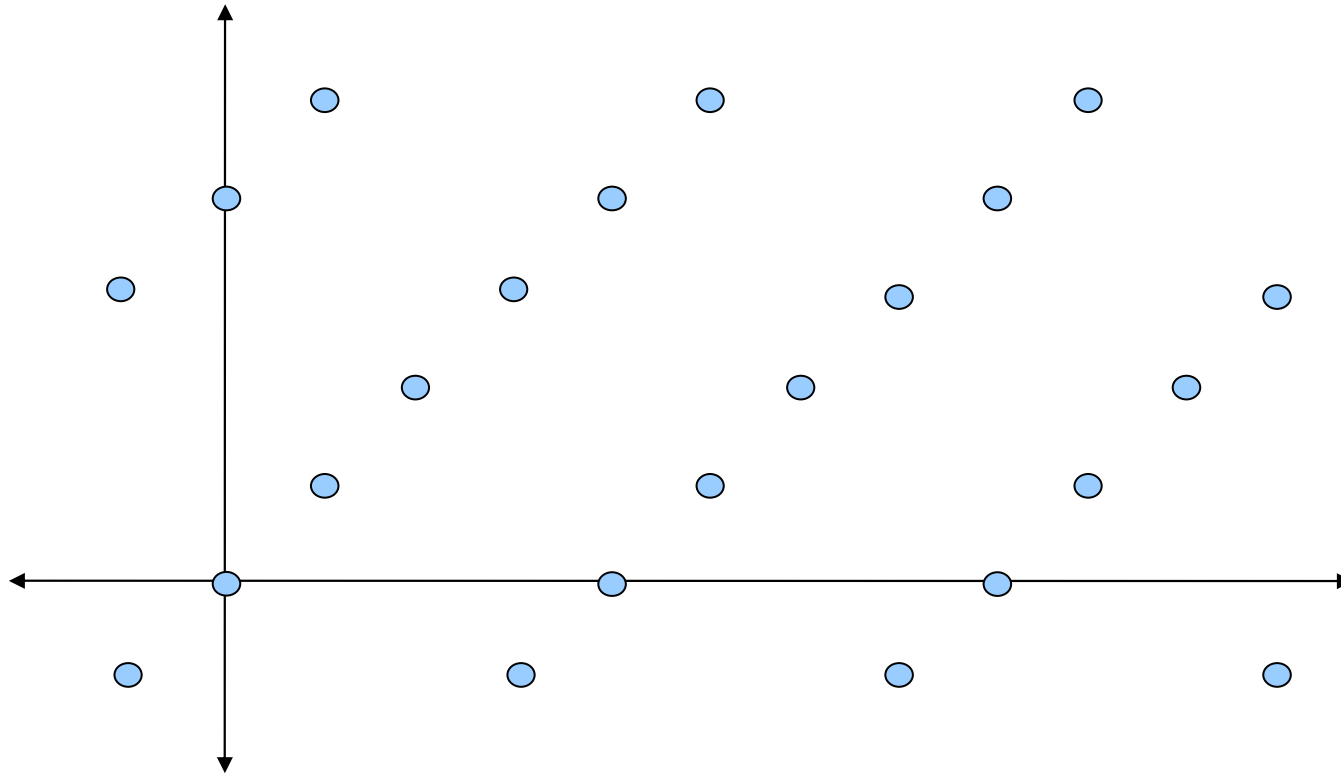
# Approximate Shortest Independent Vector Problem



Find  $n$  *pretty* short linearly independent vectors

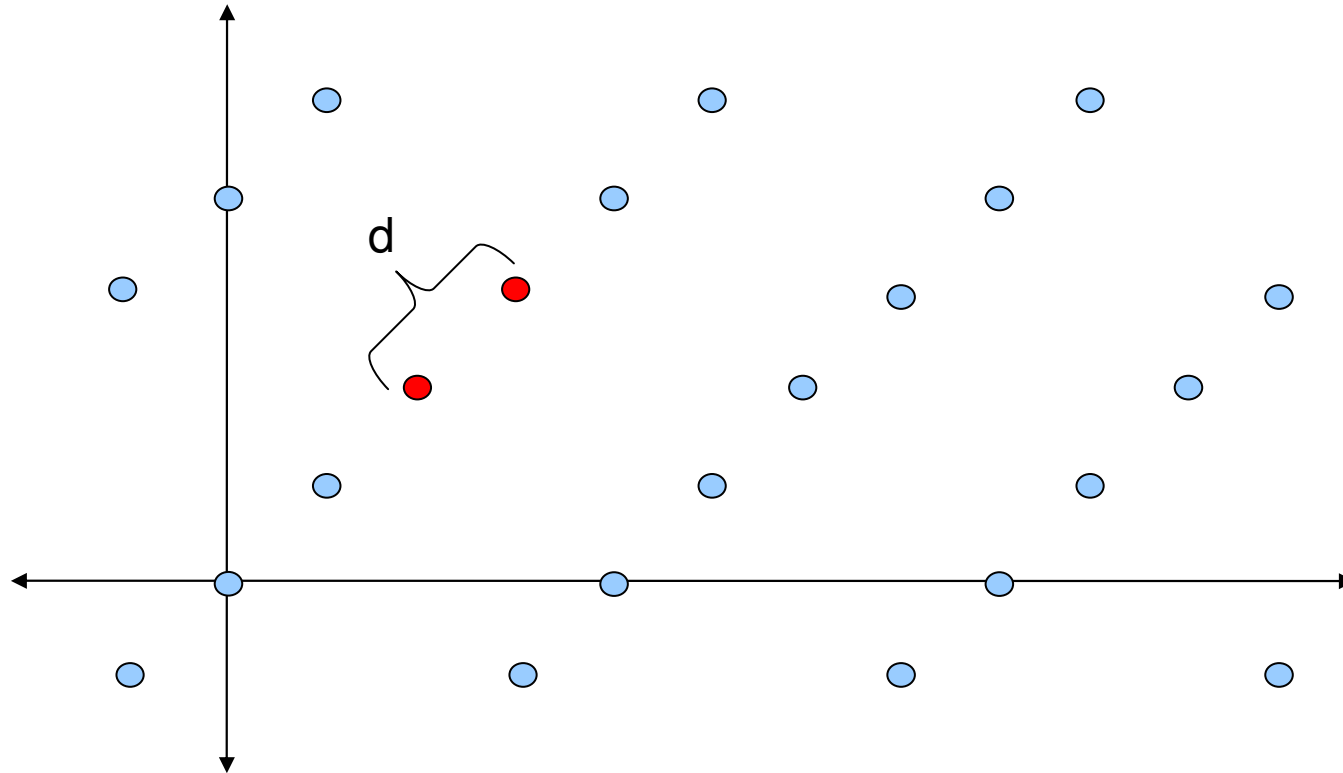


# Minimum Distance Problem (GapSVP)

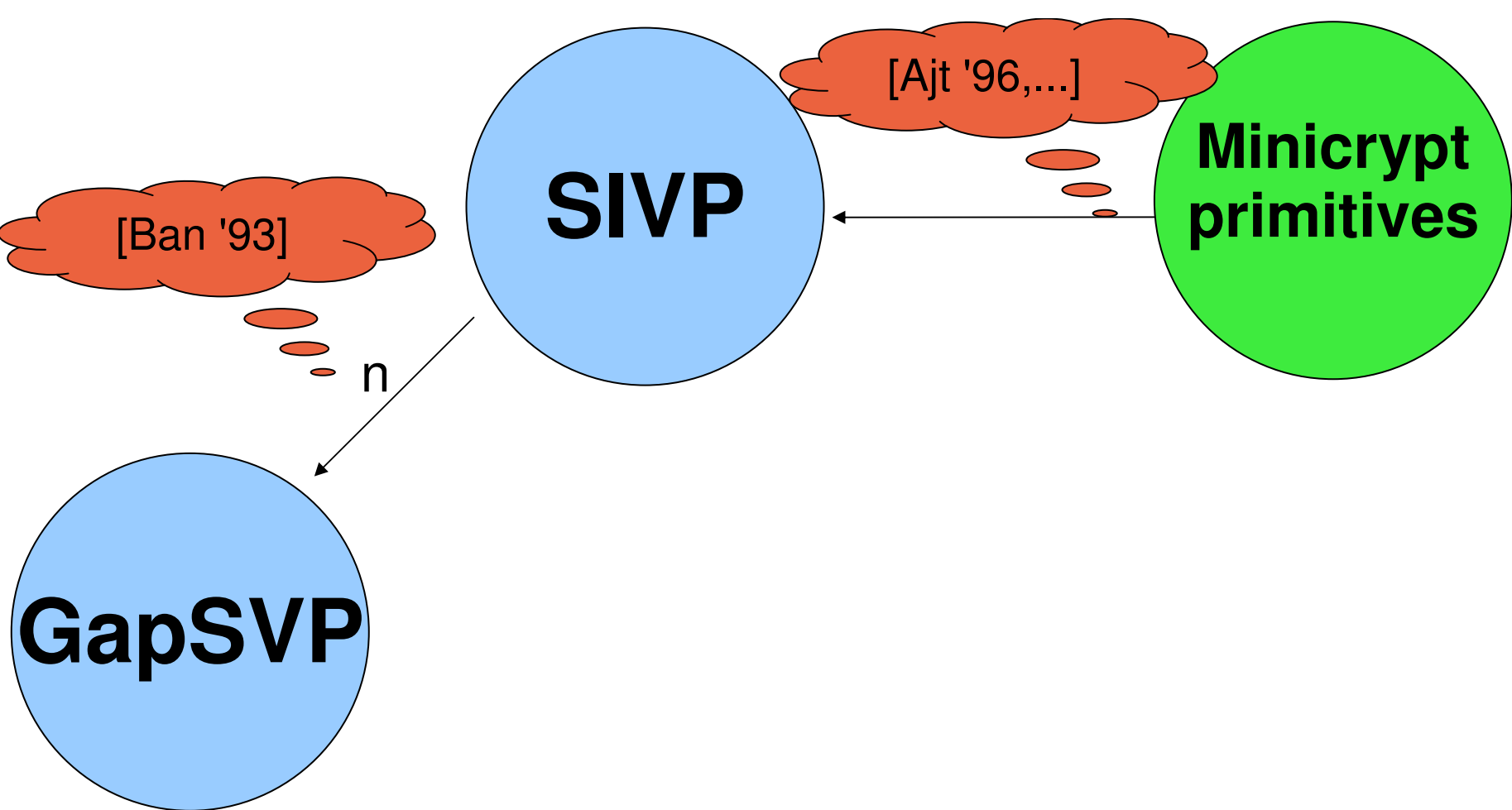


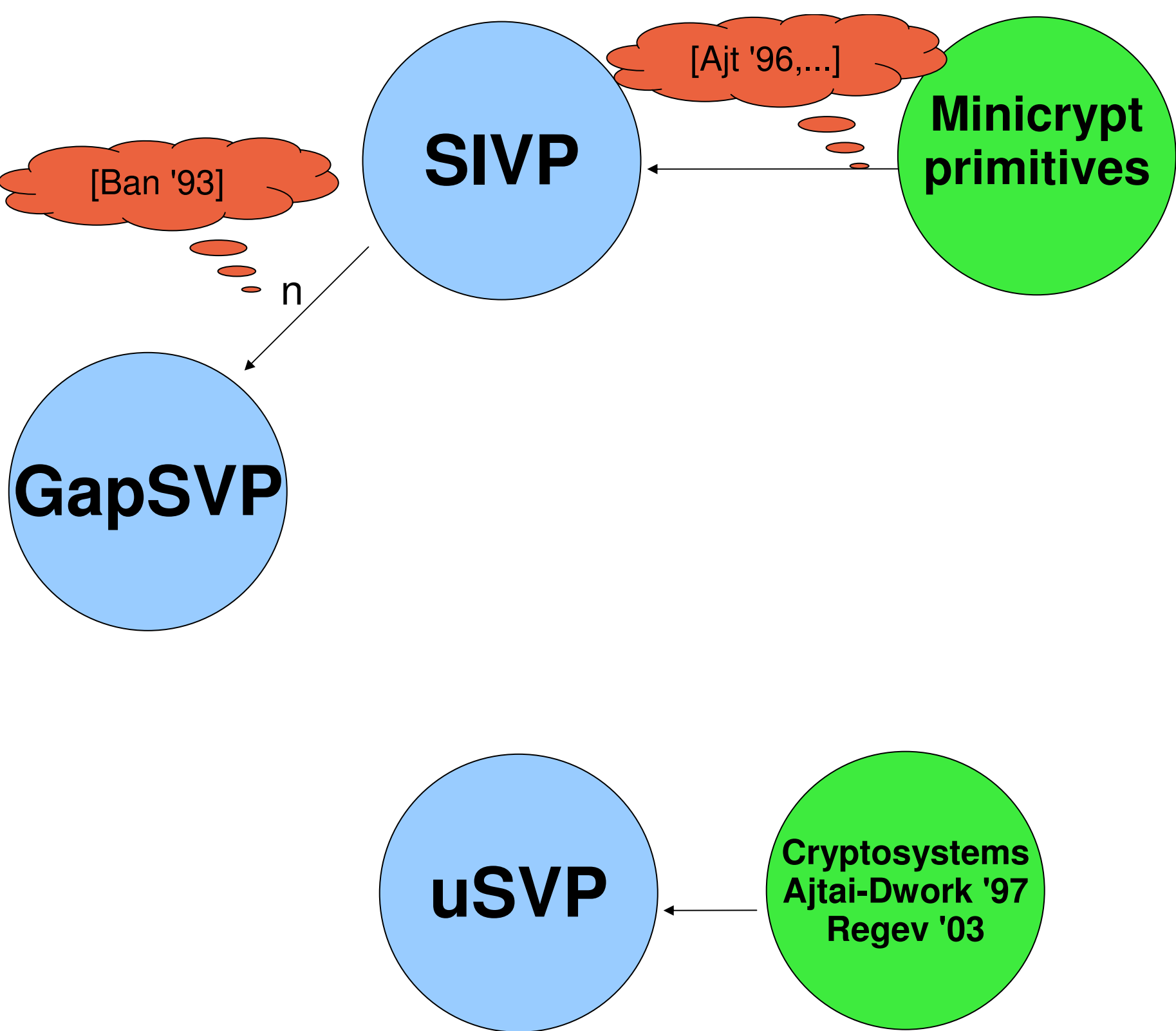
Find the minimum distance between the vectors in the lattice

# Minimum Distance Problem (GapSVP)

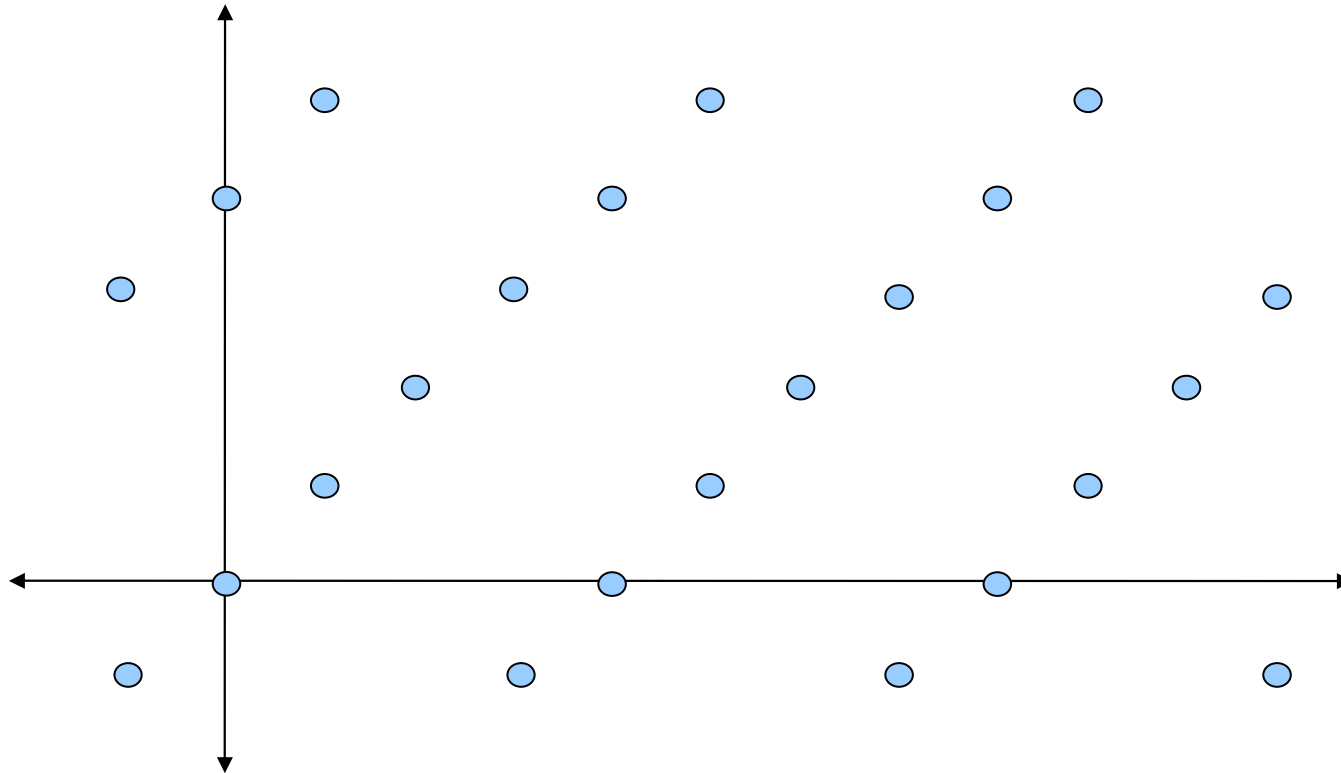


Find the minimum distance between the vectors in the lattice



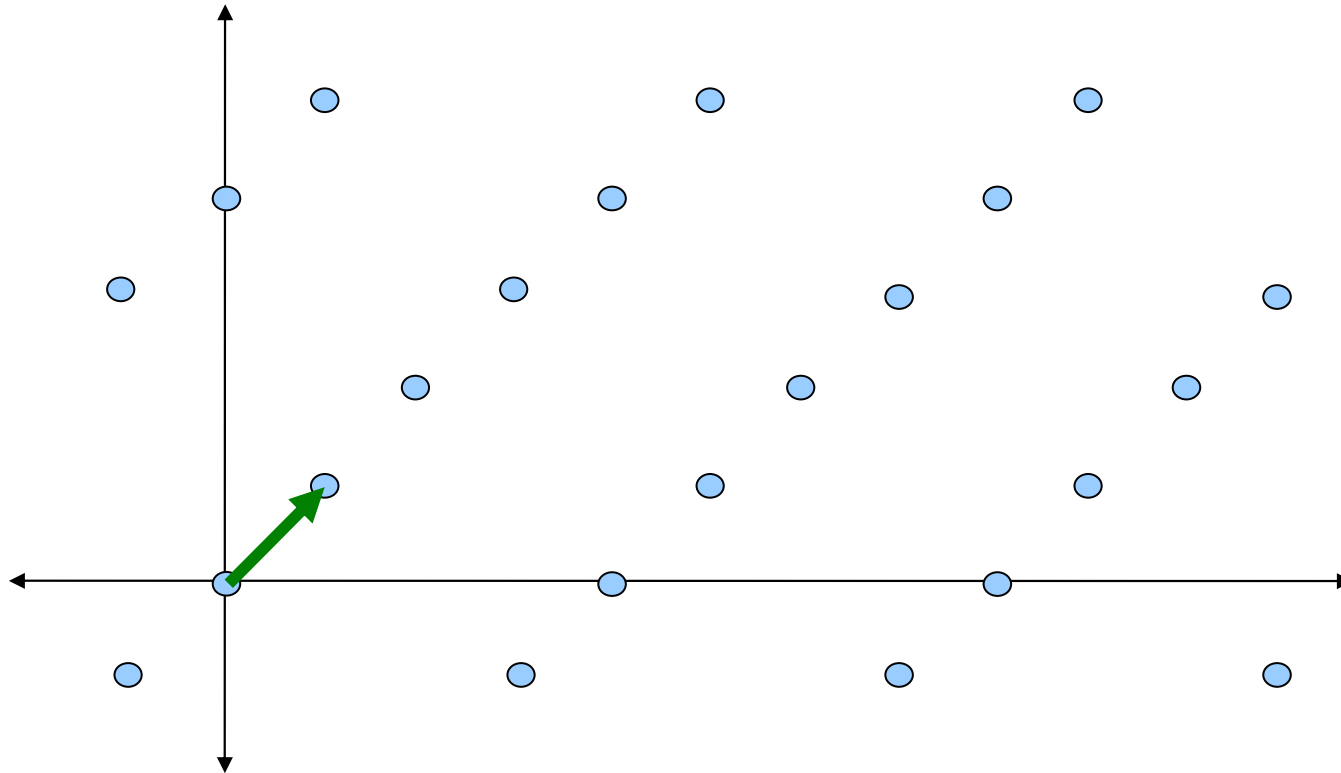


# Unique Shortest Vector Problem (uSVP)



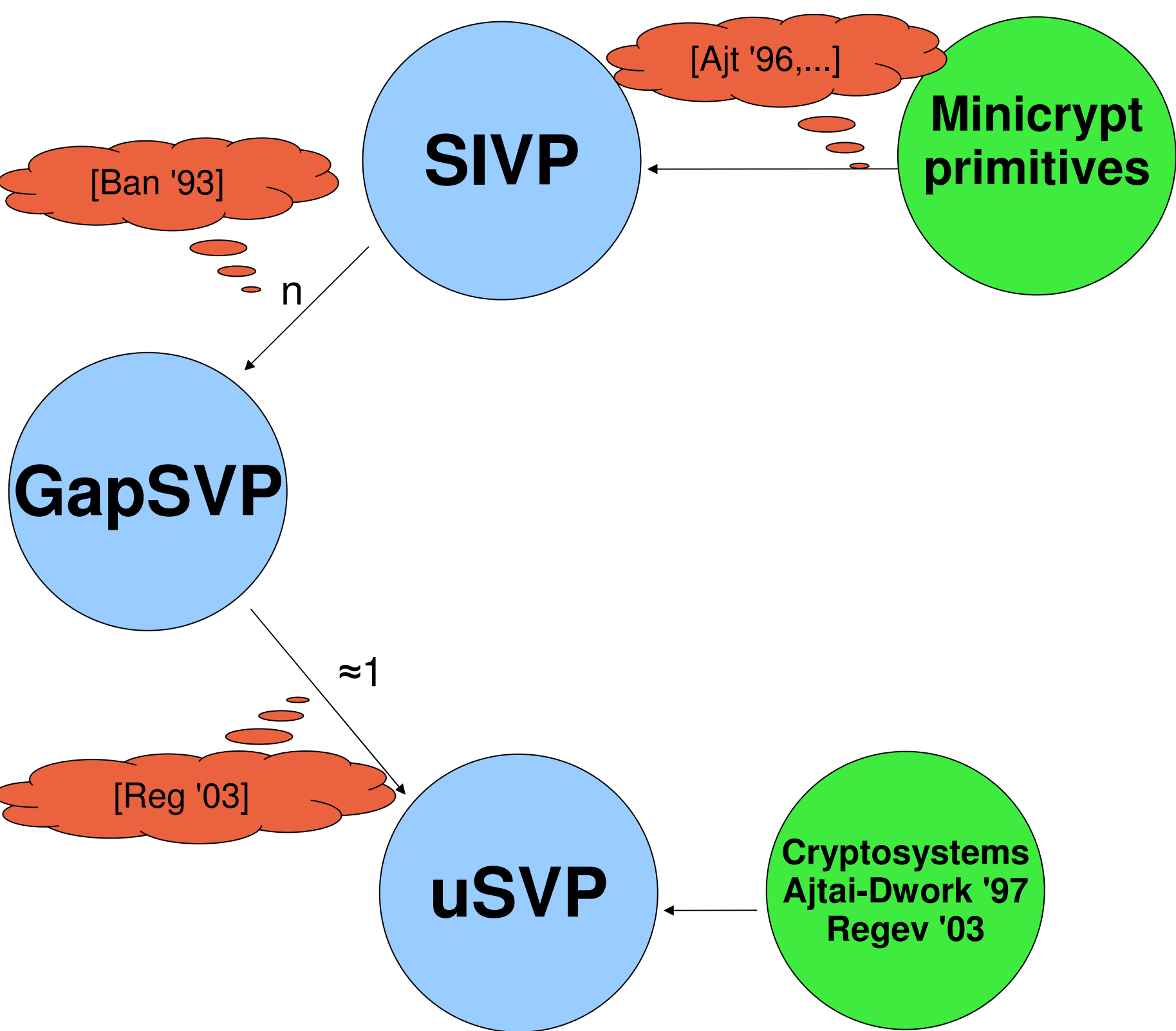
Find the shortest vector in a lattice in which the shortest vector is much smaller than the next non-parallel vector

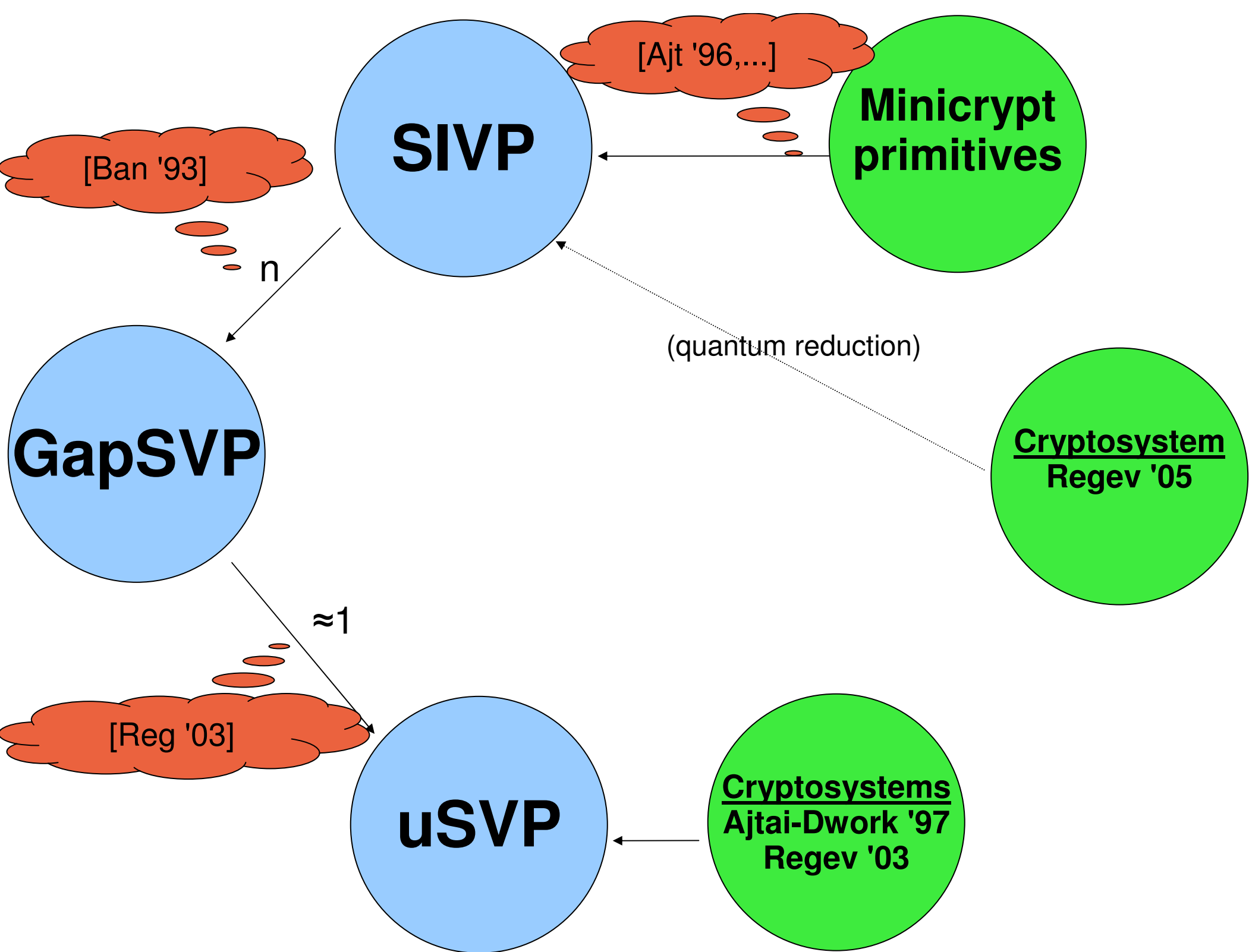
# Unique Shortest Vector Problem (uSVP)

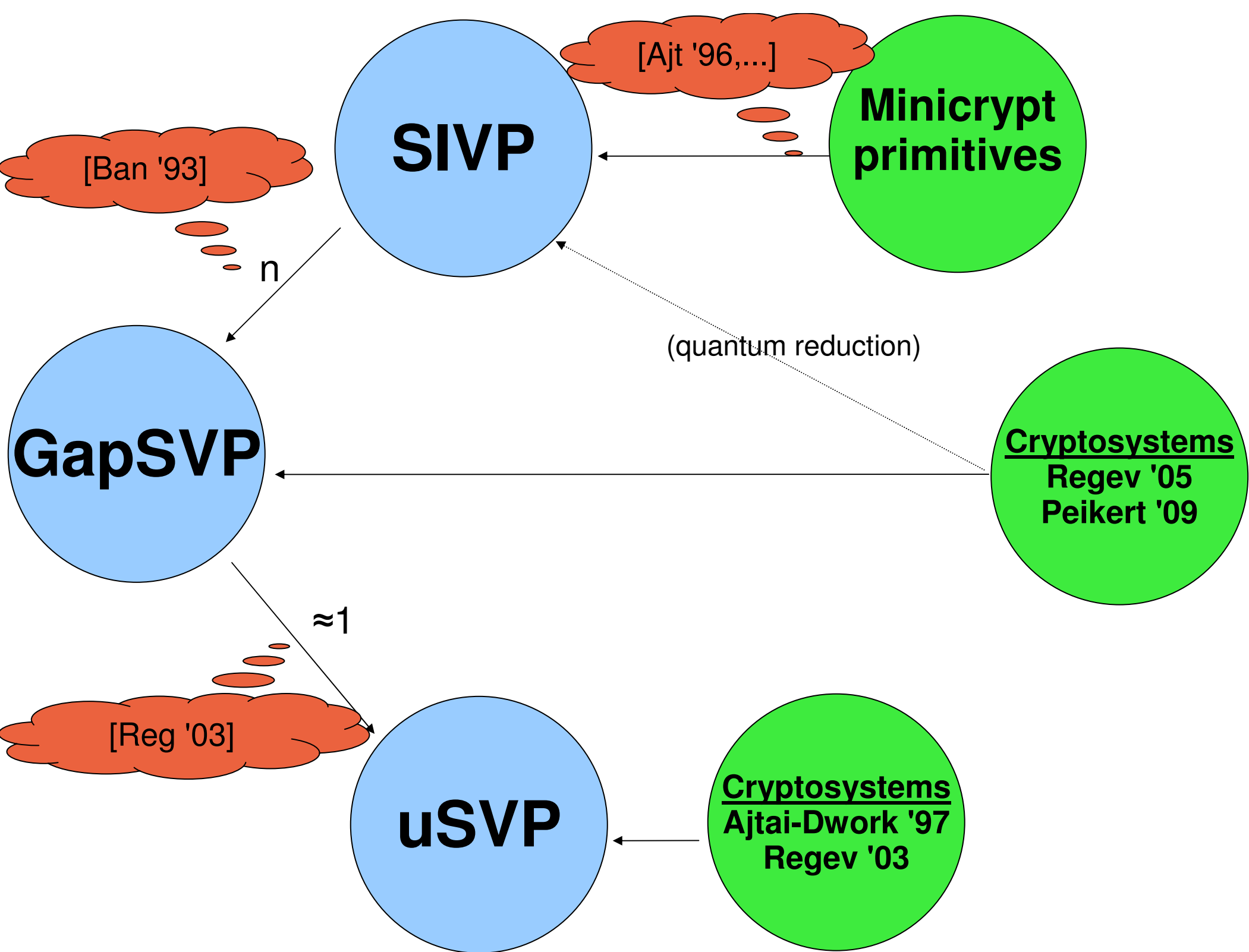


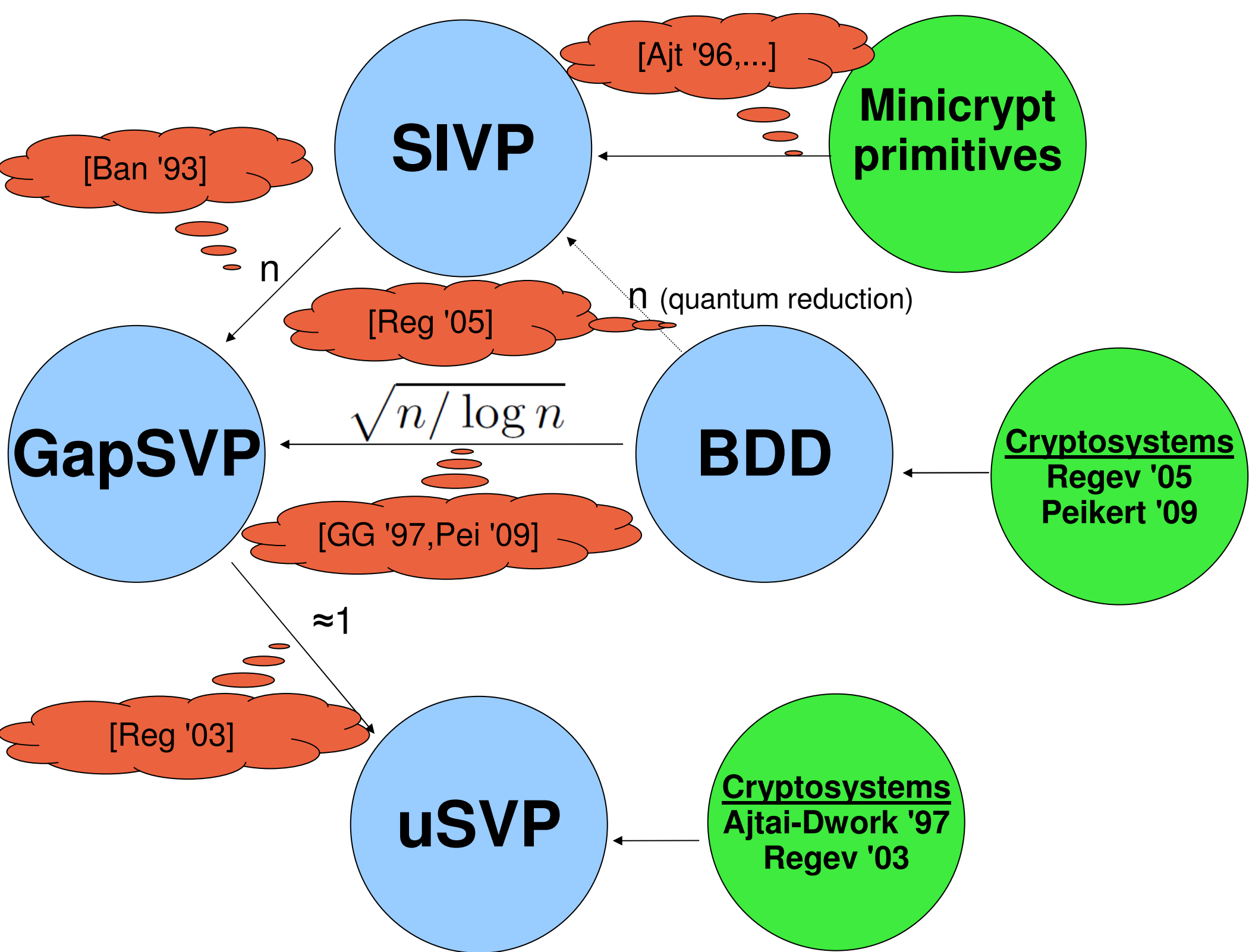
Find the shortest vector in a lattice in which the shortest vector is much smaller than the next non-parallel vector



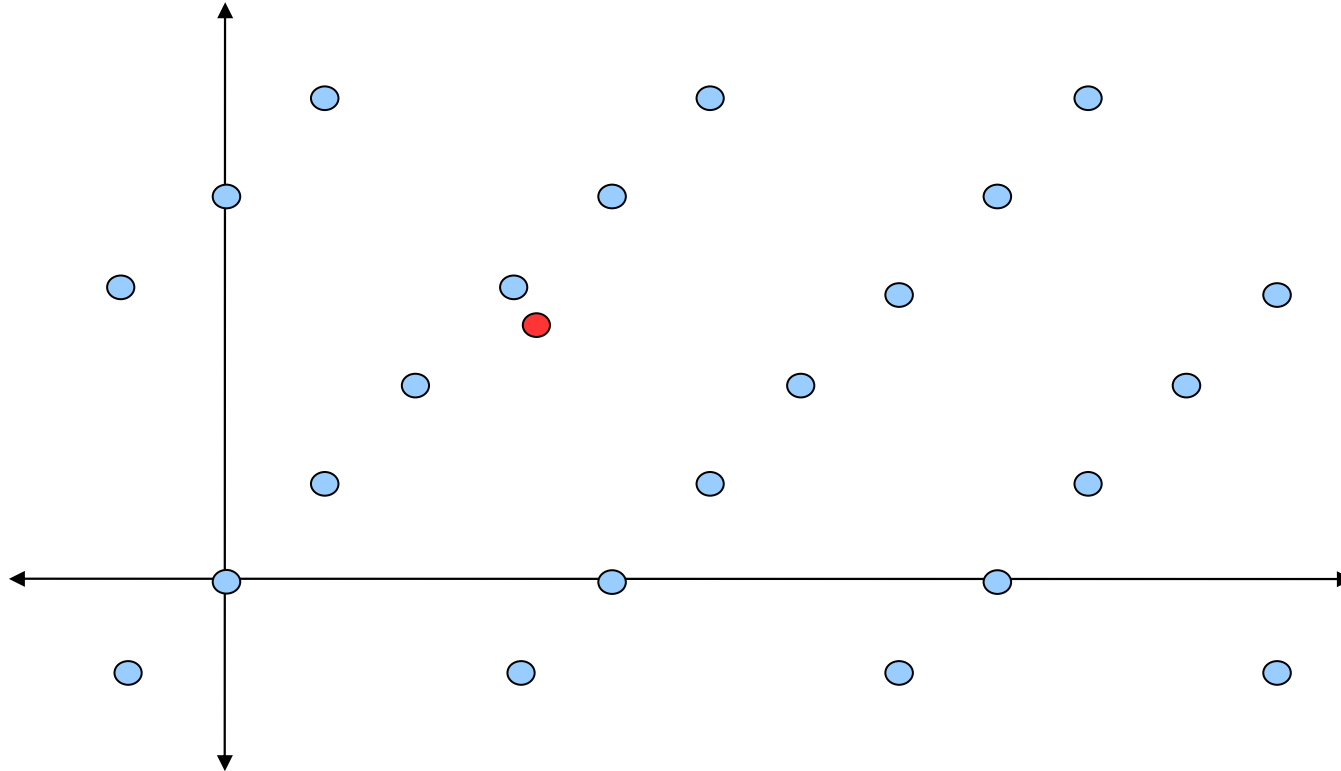




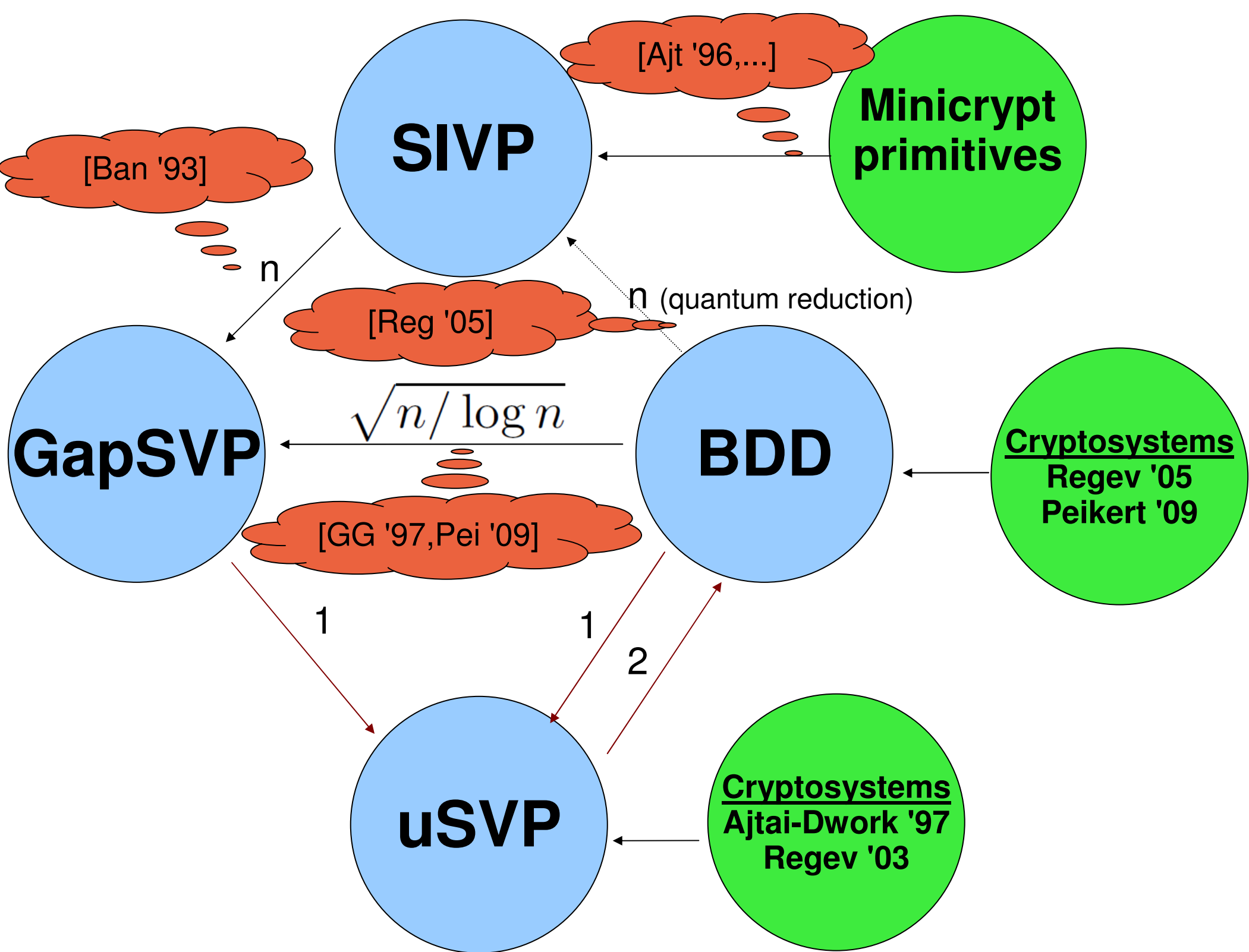


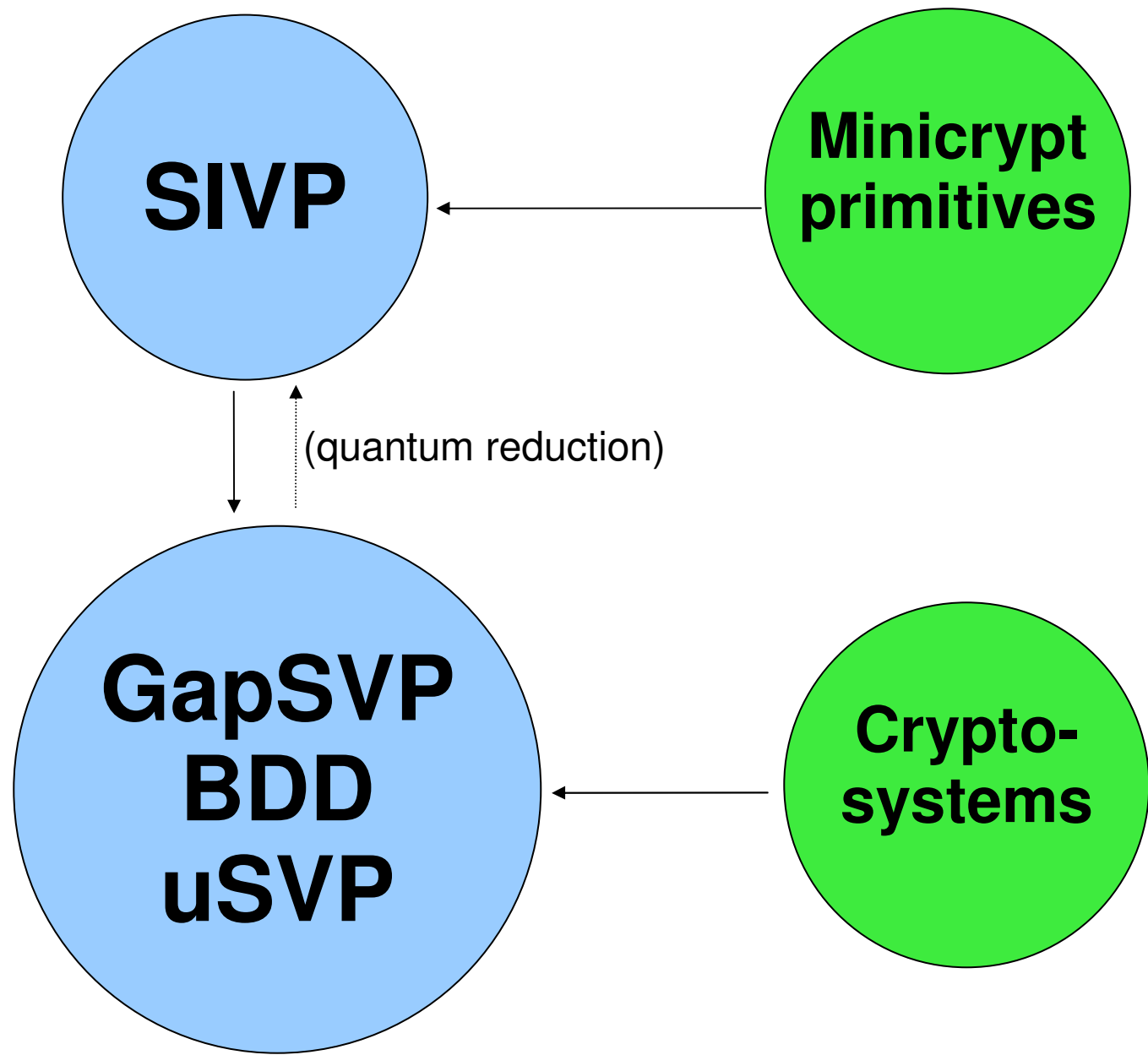


# Bounded Distance Decoding (BDD)



Given a target vector *that's close to the lattice*, find the nearest lattice vector





# Cryptosystem Hardness Assumptions

	uSVP	BDD	GapSVP	SIVP (quantum)
Ajtai-Dwork '97	$O(n^2)$	$O(n^2)$	$O(n^{2.5})$	$O(n^3)$
Regev '03	$O(n^{1.5})$	$O(n^{1.5})$	$O(n^2)$	$O(n^{2.5})$
Regev '05	-	-	-	$O(n^{1.5})$
Peikert '09	$O(n^{1.5})$	$O(n^{1.5})$	$O(n^2)$	$O(n^{2.5})$

Implications of our results



# Lattice-Based Primitives

## Minicrypt

- One-way functions [Ajt '96]
- Collision-resistant hash functions [Ajt '96, MR '07]
- Identification schemes [MV '03, Lyu '08, KTX '08]
- Signature schemes [LM '08, GPV '08]

All Based on  
GapSVP  
and SIVP

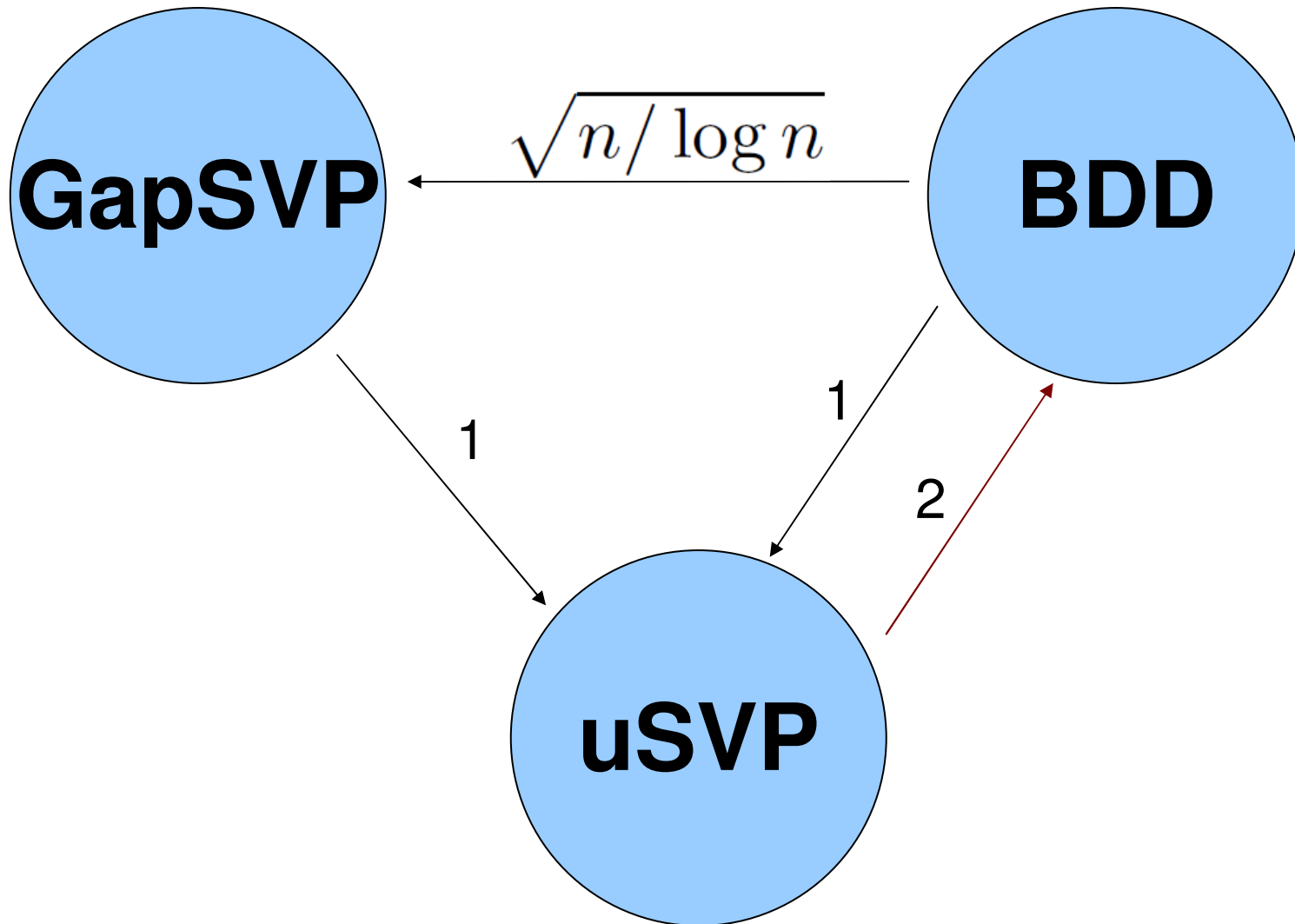
## Public-Key Cryptosystems

- [AD '97] (uSVP)
- [Reg '03] (uSVP)
- [Reg '05] (SIVP and GapSVP under quantum reductions)
- [Pei '09] (GapSVP)

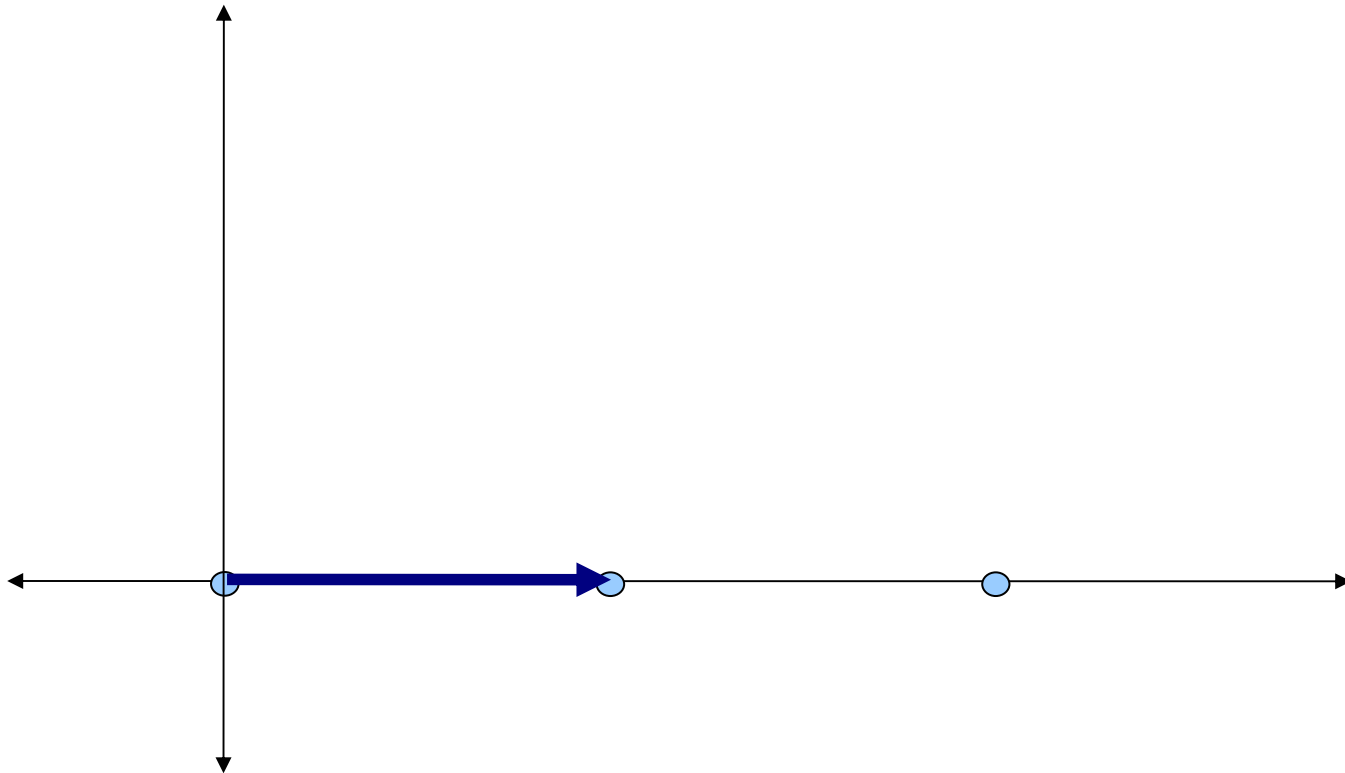
All Based on GapSVP  
and  
quantum SIVP

**Major Open Problem:  
Construct cryptosystems based on SIVP**

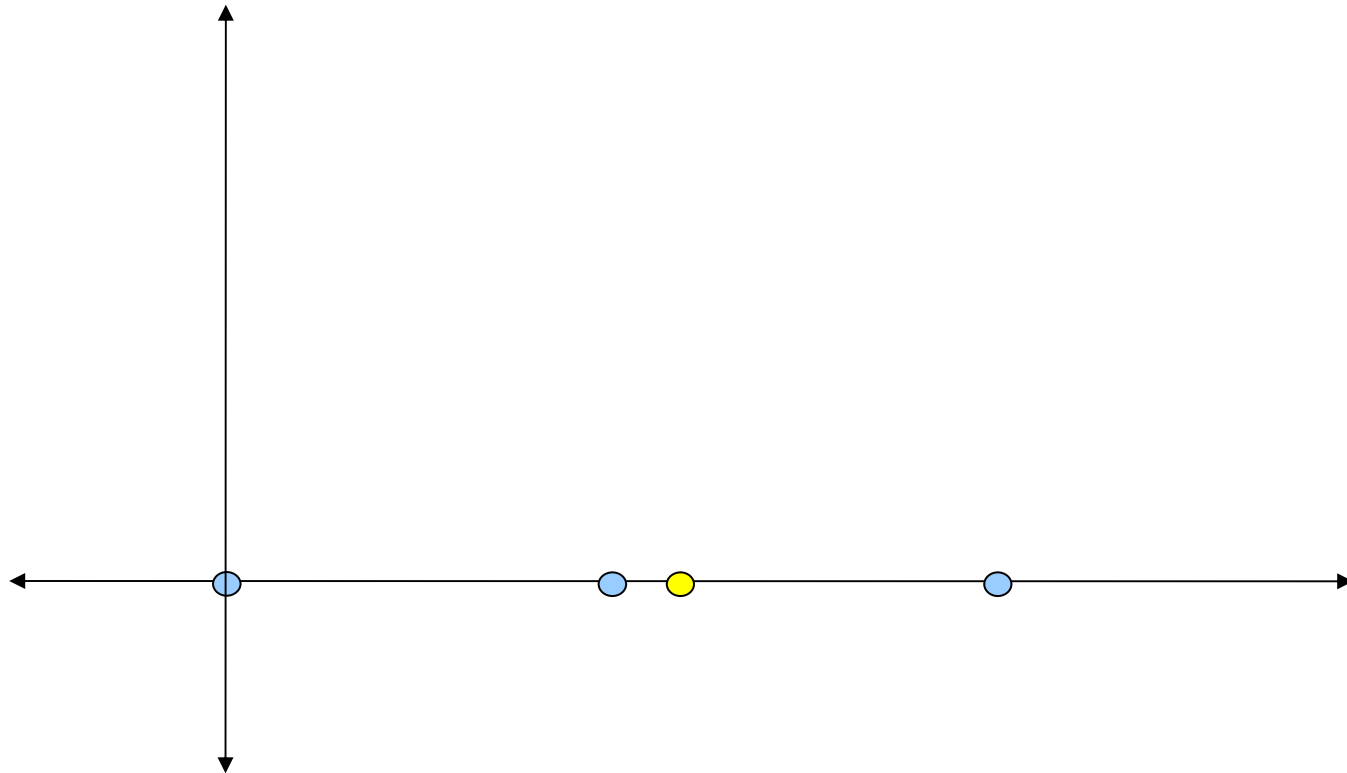
# Reductions



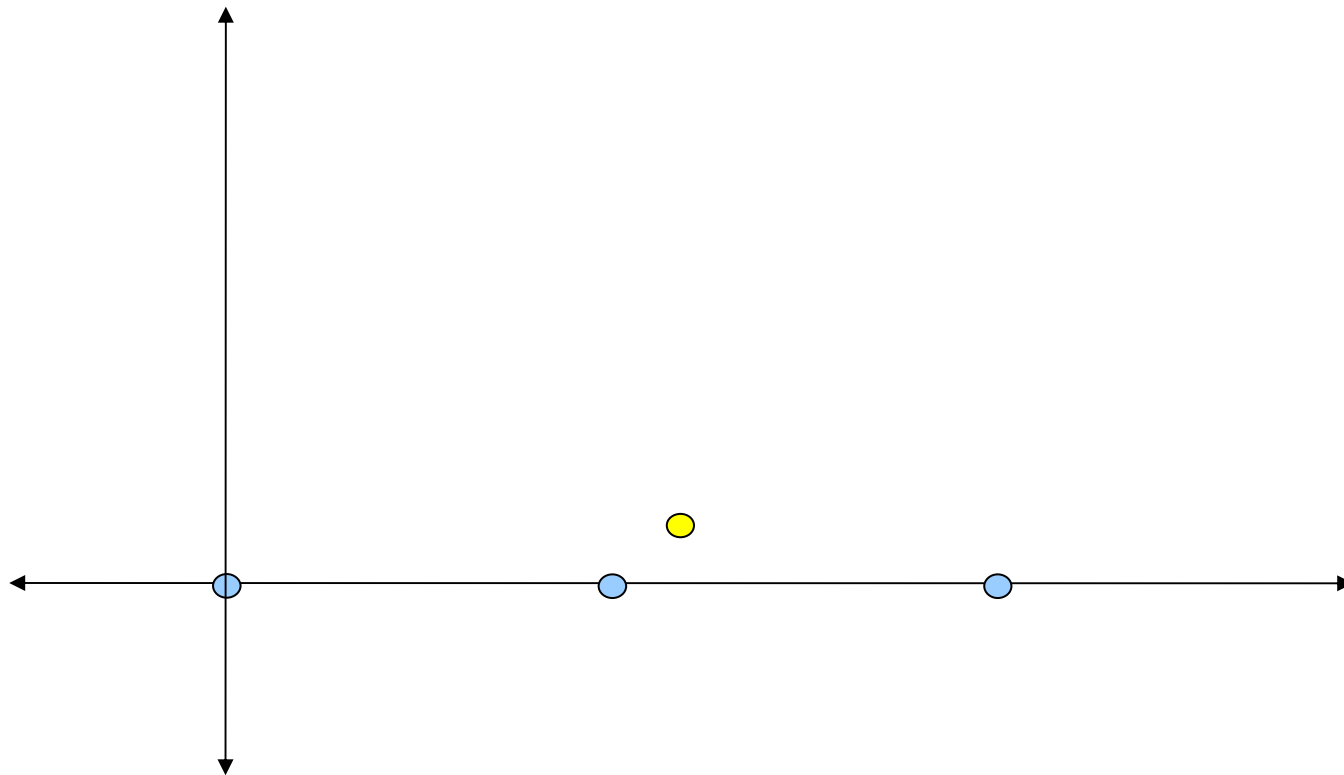
# Proof Sketch (BDD < uSVP)



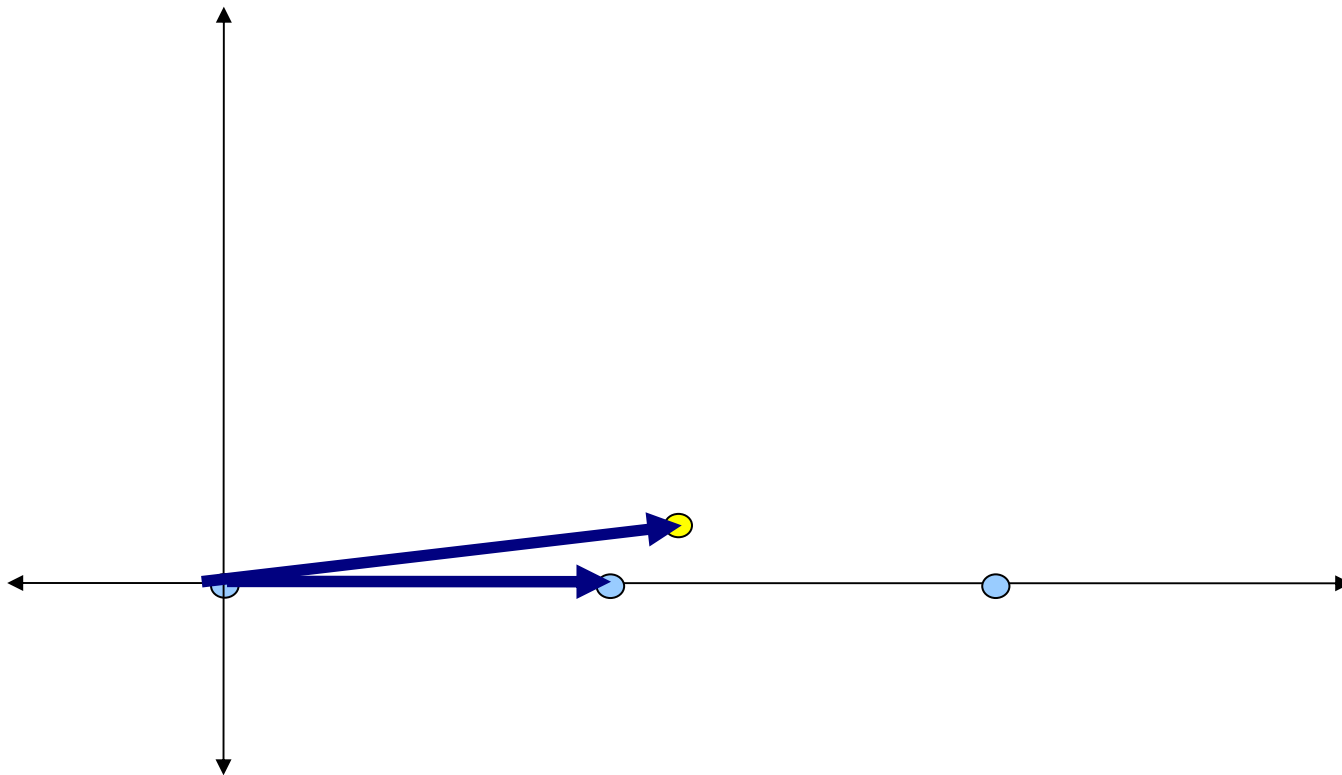
# Proof Sketch (BDD < uSVP)



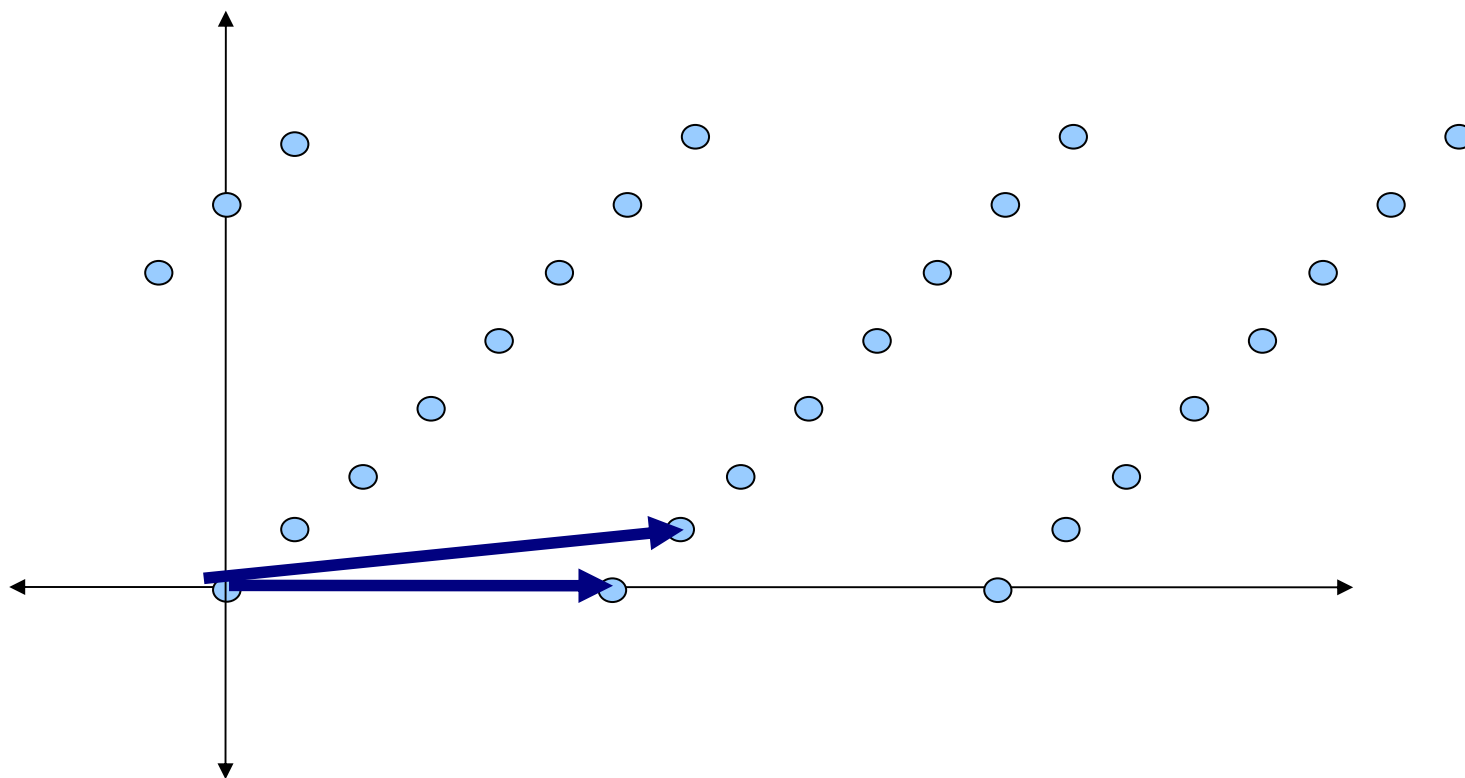
# Proof Sketch (BDD < uSVP)



# Proof Sketch (BDD < uSVP)

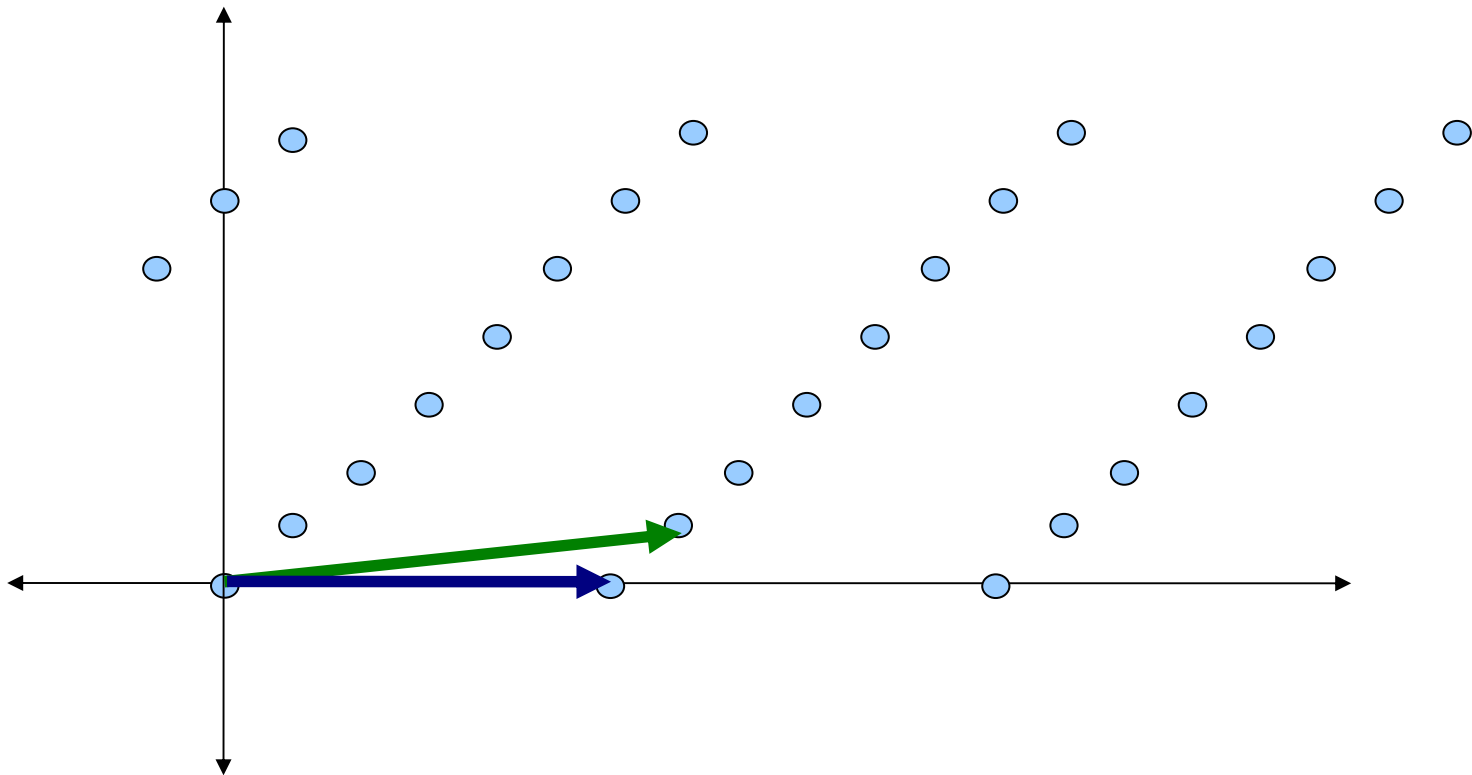


# Proof Sketch (BDD < uSVP)



# Proof Sketch (BDD $<$ uSVP)

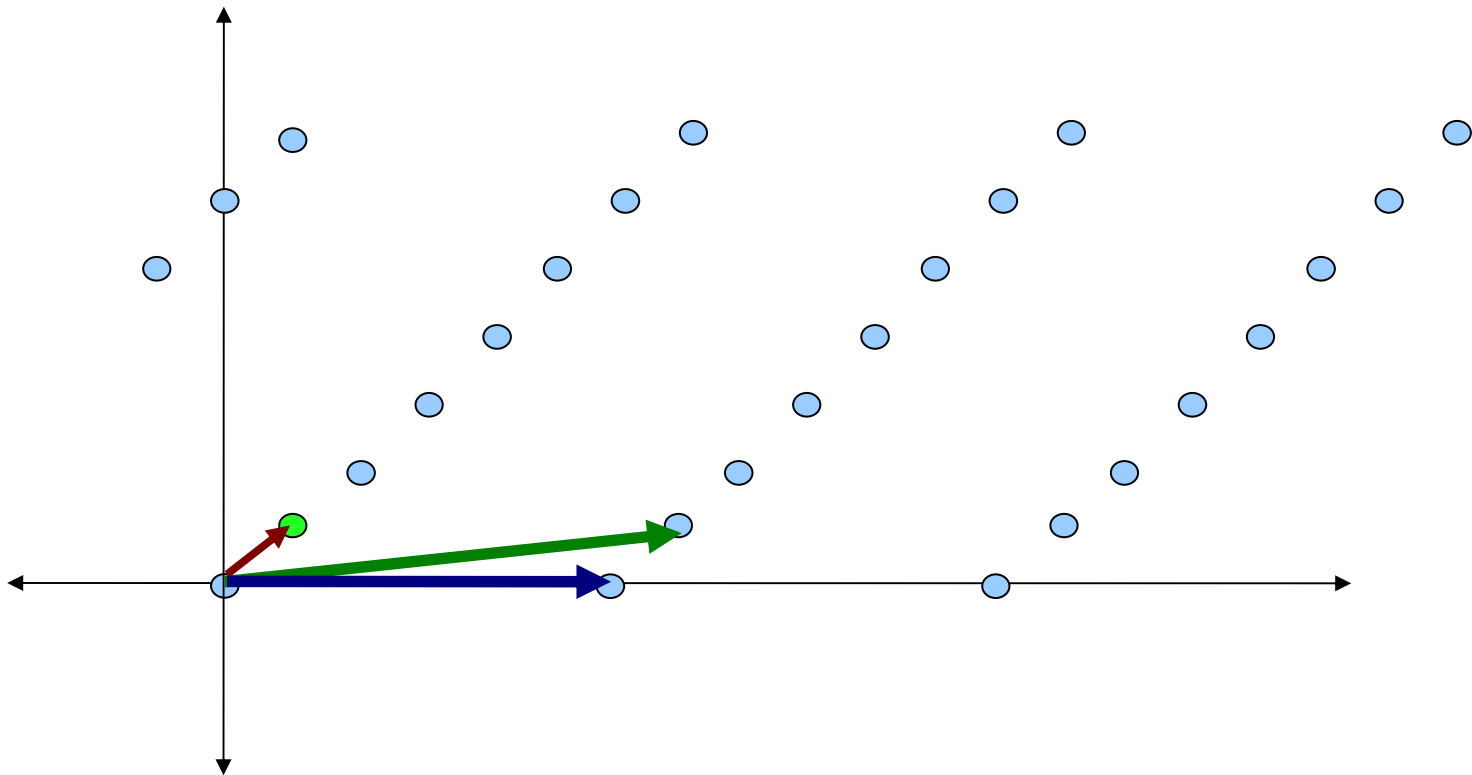
New basis vector used exactly once in constructing the unique shortest vector





# Proof Sketch (BDD $<$ uSVP)

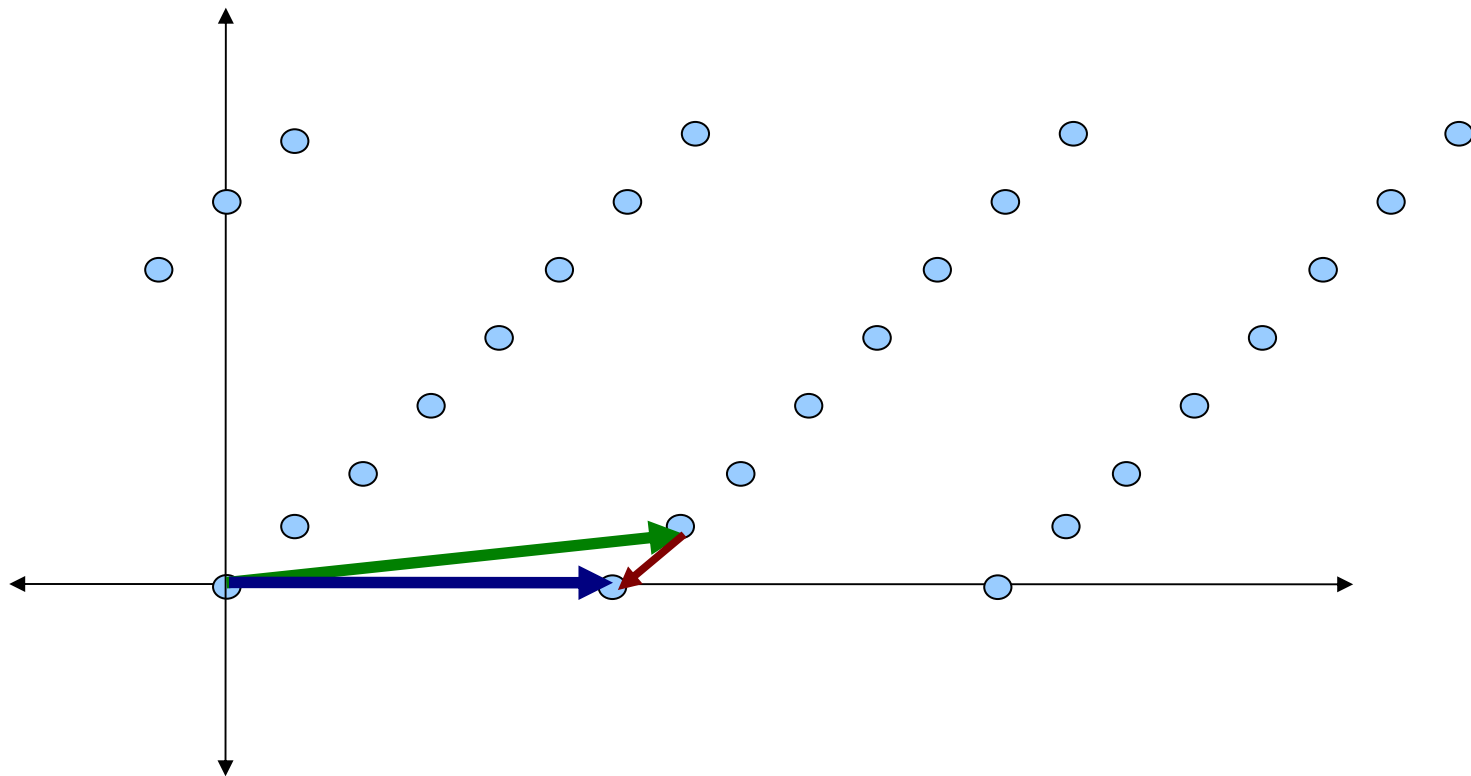
New basis vector used exactly once in constructing the unique shortest vector



# Proof Sketch (BDD $<$ uSVP)

New basis vector used exactly once in constructing the unique shortest vector

Subtracting unique shortest vector from new basis vector gives the closest point to the target.



# Open Problems

- Can we construct cryptosystems based on SIVP
  - (SVP would be even better!)
- Can the reduction  $\text{GapSVP} < \text{BDD}$  be tightened?
- Can the reduction  $\text{BDD} < \text{uSVP}$  be tightened?

**Thanks!**